



### ST. PETERSBURG INSTITUTE FOR INFORMATICS AND AUTOMATION (SPIIRAS)

EUROPEAN OFFICE OF AEROSPACE RESEARCH AND DEVELOPMENT (EOARD)

# Agent-Based Model of Information Security System: Architecture and Formal Framework for Coordinated Intelligent Agents Behavior Specification

Final Report #2

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Project Manager
Chief Scientist of SPIIRAS
Ph.D. Professor

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St. Petersburg February 2001

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The contractor will research and further develop the technology supporting an agent-based architecture for an information security system and							
a formal framework to specify a model of distributed knowledge. This research will include a comparative analysis of existing steganography & steganoanalysis techniques and the development and mathematical justification of an alternative approach to be incorporated in the							
proposed architecture.							
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#### **Preface**

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## Chapter 1. Case Study of Agent-based Information Security System: Conceptual Model, Architecture, Software Implementation and Simulation

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17C<%"/>ACA'94(O2A5(@/'%"A/C(.#"@%8(/(5'/"'A<;(3#A<'(.#"(.7"'2%"(%..#"'5(A<(8%1%C#3@%<'(/<(/;%<'= >/5%8(@#8%(#.(1\_664(D/A<(&#<&%3'5(#.("%@#'%(/''/&P5(#<(&#@37'%"(<%'F#"P5(F%"%(&#<5A8%"%84 O2%(5'/<8/"8("%@#'%(/''/&P(A5('2%("%@#'%(7</7'2#"A]%8(A<.#"@/'A#<(%..%&'E("%/CA]%8(#<(8/'/(CA<P5 '2/'( A5( '93A&/C( '#( /<9( &#@37'%"( <%'F#"P4( O2%( 37"3#5%5( #.( '2%( 3"#\$%&'( &/75%8( /( <%&%55A'9( #. %/>#"/'A#<( #.( '2%( 5/'A5./&'#"9( &#@37'%"( /<8( <%'F#"P( /''/&P( '/S#<#@A%5( /<8( /( @7C'A.#C8 &C/55A.A&/'A#<(#.('2%("%@#'%(/''/&P54(M7"A<;('2%(.A"5'(32/5%("%5%/"&2(&2/"/&'%"A5'A&5(#.(5/'A5./&'#"9 '/S#<#@A%5(#.(/''/&P5(F%"%(/</C9]%8(/<8(/("%IA%F(#.('2%(%SA5'A<; (/''/&P(/; /A<5'(5A<; C%(&#@37'%" /<8(/; /A<5'(&#@37'%"(<%'F#"P('/S#<#@A%5(F/5(3"%5%<'%84(O2%(C/''%"(A<&C78%8(/(CA5'(#.(/''/&P('%"@5E /(CA5'(#.(/''/&P(&/'%; #"A%5E(/''/&P("%57C'5(&/'%; #"A%5E(%@3A"A&/C(CA5'5(#.(/''/&P('93%5E(17C<%"/>ACA'A%5 @/'"A&%5E(/&'A#<=>/5%8('/S#<#@A%5E(5%&7"A'9(.C/F5(#"(17C<%"/>ACA'A%5('/S#<#@A%5E(/(&C/55A.A&/'A#<(#. A<'"75A#<5(>/5%8(#<('2%(5A; </'7"%5E(/<8(A<&A8%<'('/S#<#@A%5(Z'2%(5%&7"A'9('%"@A<#C#; 9(#.(&#@@#< C/<; 7/; %(.#"(5%&7"A'9(A<&A8%<'5(/"%(&#<5A8%"%8(2%"%[4(?'(F/5(&#<&C78%8('2/'('2%(A<&A8%~'('/S#<#@A%5 /"%( '2%( @#5'( 3"#@A<%<'( .#"( ! "#\$%&'( ; #/C5E( >7'( '2%9( "%U7A"%( /( 8%%3%"( /</C95A5( FA'2( %@32/5A5( #< "%@#'%( /&'A#<54( O2A5( 3/"'( #.( '2%( "%5%/"&2( "%57C'%8( A<( '2%( 8%1%C#3@%<'( #.( '2%( &C/55A.A&/'A#<( #. 5'/<8/"8("%@#'%(/''/&P54(?<(''2A5(&C/55A.A&/'A#<(%/&2(/'''/&P(53%&A.A%5(>9('2%(A<'%<8%8(%...%&'E(37"3#5% #.('2%(/''/&PE(&#<8A'A#<5(3"#1A8A<;(.#"('2%(%..%&'(/&2A%1%@%<'E(%SA5'%<&%(#.(/(.%%8>/&P(FA'2('2% /''/&P%8(#>\$%&'E(C/9#7'(#.('2%(57>\$%&'(#.(/''/&P(&#<&%"<A<;('2%(/''/&P%8(#>\$%&'E(C/9%"(#.(5'/<8/"8 ?6HaH6?(@#8%C(#<(F2A&2('2%(%..%&'(A5(&/""A%8(#7'E '2%(#>\$%&'('2/'(A5('2%(/''/&P('/"; %'4(\V%5A8%5E '93A&/C(5&2%@%5(#.('2%("%@#'%(/''/&P5(@/33%8('#('2%(8%1%C#3%8(&C/55A.A&/'A#<(F%"%(A<'"#87&%84(?' F/5(8%5&"A>%8(5AS(; %<%"A&(5&2%@%5(#.('2%("%@#'%(/''/&P5('2/'(/"%(/5(.#CC#F5B

- /</C95A5(#.('2%(<%'F#"P('"/..A&(Z#"(CA5'%<A<; (#.(/(8/'/(CA<P(>9(@%/<5(#.(53%&A/C('##C5(o 5<A..%"5[E
- <%'F#"P(5&/<<A<; E
- 57>5'A'7'A#<( #.( '2%( '"75'%8( #>\$%&'( #.( '2%( <%'F#"P( /<8( '"/<5@A55A#<( #<( 8/'/( CA<P5( #.( '2% @%55/; %5(."#@(A'5(</@%(FA'2(/33"#3"A/'A#<(#.(A'5(/&&%55("A; 2'5E
- A@3C/<'/'A#<(#.('2%(./C5%(#>\$%&'(#.('2%(<%'F#"PE
- 8%<A/C(#.(5%"I A&%E
- "%@#'%(%S%&7'A#<(#.(/33CA&/'A#<54

O2%(3"#3#5%8(/"&2A'%&'7"%(&#@3"A5%5(/(<7@>%"(#.(53%&A/CA]%8(&##3%"/'A<; (5%&7"A'9(/; %<'54O2%(5%'(#.(/; %<'5(A<&C78%5

- Intrusion detection agents("%53#<5A>C%(.#"(8%'%&'A#<(A<'"75A#<(#<('2%(>/5A5(#.(A<.#"@/'A#< "%&%A1%8(."#@(#'2%"(/;%<'5(/>#7'(3#55A>C%(/<#@/CA%5E(5753A&A#75(>%2/1A#"(#.(75%"5E(/''%@3'5#.(<#<=/7'2#"A]%8(/&&%55(%'&4
- *Access control agents*(3"#1A8A<; (/<(/&&%55('#('2%(A<.#"@/'A#<("%5#7"&%5(A<(/&&#"8/<&%(FA'2 '2%(75%"5n(/7'2#"A'94
- *Identification and authentication agents*( "%53#<5A>C%( .#"( A8%<'A.A&/'A#<( #.( A<.#"@/'A#< 5#7"&%5(/<8(&#<.A"@/'A#<(#.('2%A"(/7'2%<'A&A'94

O2%5%(/; %<'5( F%"%( &#<5A8%"%8( /5( >/5A&( #<%5( A<( '2%( .7'7"%( 8%I %C#3@%<'( #.( '2%( @7C'A=/; %<' 1\_66(/"&2A'%&'7"%(/<8(A'5(5#.'F/"%(A@3C%@%<'/'A#<4(?<(; %<%"/C(&/5\%(A'(A5(@/P%(5%<5%('#(%S'%<8('2%/%&17%('2"#7; 2(A<&C78A<; (5#.'F/"%(/; %<'(#.(#'2\%"(PA<8('2/'(/"%

- Attack suppression agents("%/CA]A<; (p3"#5%&7'A#<q(/<8(<%7'"/CA]/'A#<(#.('2%(/''/&PA<; (3"#; "/@5 ZA<'"78%"5[E
- Agents of damage assess( /<8( A<.#"@/'A#<( A<'%; "A'9( "%8#1%"9( "%53#<5A>C%( .#"( A<.#"@/'A#< A<'%; "A'9(%<57"A<; E
- Cryptography and steganography( 3"#'%&'A#<( /; %<'5( 75%8( .#"( &"%/'A#<( #.( 5/.%=%S&2/<; % &2/<<%\5/>%'F\%<(<%'F\#"P(<\#8\5(\72\#5'\5E(\5\"1\8"\5)[E

• Learning agents( "%53#<5A>C%( .#"( /8/3'/\h#<( #.( 1\_66( /; %<'5( '#( &#@37'\%"( <\"F#"P "\&#<.A; 7"/\h#<(/<8(/''\&P5(#.(<\%F('93\%54

?<(5#%(&/5%5(A'(&/<(>%("%/5#</>C%('#(A<&C78%(A<(/"&2A'%&'7"%(/(5#=&/CC%8(@%'/=/; %<'"%53#<5A>C%(.#"(@/</; A<; ('2%(A<.#"@/'A#<(5%&7"A'9(3"#&%55%5E(3"#1A8A<; (.#"(&##3%"/'%8(>%2/1A#"(#.'2%(5%&7"A'9(/; %<'5(/<8("%/CA]A<; ('2%("%U7A"%8(C%1%C(#.(; %<%"/C(5%&7"A'9(A<(/&&#"8/<&%(FA'2(/(; C#>/C&"A'"A#<4

O2%(57>\$%&'5(#.('2%(Interim Report #2(f?<'J%3)-g(F%'%('2%("%57C'5(#.("%5%/"&2(#<('2%.#CC#FA<; ('/5P5B(Z\*[(8%1%C#3@%<'(#.('2%(@#8%C5(#.(8A5'"A>7'%8(&#@@#<(/<8(53%&A.A&(P<#FC%8; %(#./;%<'5E(Z-[(#>\$%&'=#"A%<'%8(8%5A;<(#.('2%(5#.'F/"%('##C(3"#'#'93%(#.(/;%<'=>/5%8(1\_66(/<8(Z0[8%1%C#3@%<'(#.('2%(3/"'A&7C/"(5#.'F/"%(&#@3#<%<'54

 $: (.7"'2\%''(/''\%<'A\#<(F/5(3/A8('\#('2\%(8\%1\%\#3@\%<'(\#. '2\%(8\#@/A< ontology /5(/(.\#''@/C."'/@\%F\#"P(.\#''(5"'7\&'7"A<;(8A5'"A>7'\%8(P<\#FC\%8; \%(>/5\M(O2\%(3"\#3\#5\%8( ontology-based framework .\#''(8\1\\%\#3@\%<'(/<8("\%3"\%5\%<'/'A\#<(\#.('2\\%(1__66(8A5'"A>7'\\%8(P<\#FC\\%8; \\(F/5(A<'\%<8\\%8('\#(3"\#1A8\\%(.\#")\@/A<'\%<'\%\(\#1A\)(\framework .\#''\)(\framework .\#'\)(\framework .\#\)(\framework .\#$ 

The most significant "dimensions" of the entire ontology of the security concepts and tasks F%"%( "%3"%5%<'%8( >9( '2%( p5%&7"A'9( /; %<'5( 8#@/A<qE( p@%55/; %=%1 %<'=/''/&P( 3"#&%55A<; qE( p5A<; C%/''/&P( '93%5qE(p5A; <A.A&/<'(%1 %<'5qE(p5%&7"A'9(/; %<'5o(/'''/&P(&C/55%5o(5A; <A.A&/<'(%1 %<'5qE(p5%&7"A'9/; %<'5n(#3%"/'A#<5q(/<8(p5&%</"A#5(#.(/'''/&P5o(5%&7"A'9(/; %<'(&##3%"/'A#<q(&#@3#<%<'54

The "security agents' domain" ontology 5%'5( 3/"'A&7C/"( .7<&'A#</CA'9( /<8( /"%/( #. "%53#<5A>ACA'9( '#(%/&2( 3/"'A&7C/"( 5%&7"A'9( /; %<'4(?'( ; #1%"<5( '2%( %<'A"%( 8%1 %:#3@%<'( &9&C%( #. ( '2% 1\_66(&#@3#<%<'54

The "message-event-attack processing" ontology 53%&A.A%5(3"#&%55A<; (#.(@%55/; %5(."#@C#F%"('#(733%"(C%1%C5(#.(;%<%"/CA]/'A#<(/<8(&C/55A.A&/'A#<4(O2%(3/"'A/CC9(#"8%"%8(5%'(#.(>/5A&&#<&%3'5(A<&C78%5(@%55/;%('"/..A&E(%1%<'E(5A;<A.A&/<'(%1%<'E(&"A'A&/C(%1%<'E(/78A'("%&#"8E(3/''%"<E"7C%5E(/&"A1A'9E(5&%</"A#(#.(/&"A1A'9E(&#<<%&"A#<E(5A@3C%(Z\5A<; C%=32/5%[(/''/&PE(&#@3#5A'%(Z\07C'A=32/5%)[(/''/&P5E(#"8%"5(&C/55%5(#.2/5%))[(/''/&P5E(#"8%"5(&C/55%5(#.2/5%))](/''/&P5E(#"8%"5(AC/55%))](/''/&P5E(#"8%"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/55%))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#"8\*"5(AC/5\*\*))](/''/&P5E(#\*"6\*\*))](/''/&P5E(#\*"6\*\*))](/''/&P5E(#\*"6\*\*))](/''/&P5E(#\*"6\*\*))](/''/&P5E(#\*"6\*\*))](/''/&P5E(#\*"6\*\*))](/''/&P5E(#

The "significant events" ontology 8%.A<%5(/(2A%"/"&29(#.(%1%<'(3"#&%55A<; (C%1%C5E(/33"#3"A/'% <%'F#"P(3"#'#&#C5(/<8(5%"1A&%5E(#3%"/'A<; (595'%@(&#@@/<85(/<8(/33CA&/'A#<54(?'(3C/95(/<(A@3#"'/<' "#C%(A<("%53%&'('#(@7C'A=C%1%C(8/'/(3"#&%55A<; (3%".#"@%8(>9(5%&7"A'9(/; %<'54

The "security agents – attack classes – significant events" ontology 53%&A.A%5('2%(/"%/(#. "%53#<5A>ACA'9(#.(%/&2(3/"'A&7C/"(A<'"75A#<(8%'%&'A#<(/; %<'4(?'(8%.A<%5(/(2A%"/"&29(#.(%1 %<'(3"#&%55A<; C%1 %C5(@/33%8(A<'#(?M: 5n(&C/55%54

The "scenarios of attacks — security agent cooperation" ontology 53%&A.A%5(&#@3#5A'%(2#5'= >/5%8(/<8(<%'F#"P=>/5%8(8A5'"A>7'%8(/''/&P5(@/33%8(A<'#(5&%</"A#5(#.(A<'%"/&'A#<5(/<8(&##3%"/'A#<#.('2%(5%&7"A'9(/;%<'54(O2%(8%5&"A>%8(p8A@%<5A#<5q(#.('2%(%<'A"%(#<'#C#;9("%.C%&'('2%(/7'2#"5(1A%F(/'/%#@3C%SA'9(#.('2%(1\_\_66(@#8%C(/<8("%53%&'A1%(5#.'F/"%(3"#'#'93%('2/'(FACC(@/P%(A'(3#55A>C%('#%S3C#"%(/<8('#(7<8%"5'/<8(/81/<'/;%5E(8A5/81/<'/;%5(/<8(3"#>C%@5(#.(75A<;(A<'%CCA;%<'(/;%<''%&2<#C#;9(.#"(A<'%;"/C(/557"/<&%(#.('2%(<%'F#"P(5%&7"A'94

 $O2\%(Final\ Report\ \#1(fYA<J\%3)*g(57@@/"A]\%8('2\%("\%57C'5(\#.('2\%("\%5\%'"\&2(87"A<; ('2\%('2A"8(/<8\%7"'2(m7/"'\%'5(\#.('2\%("\%5\%/"&24(M7"A<; ('2A5(3\%"A\#8('2\%(.\#Ct\#FA<; ('/5P5(F\%"\%5\#C\#8B)$ 

- \*[(M%1%#3@%<'(#.('2%(@#8%5(#.(8A5'"A>7'%8(&#@@#<(/<8(53%&A.A&(P<#FC%8; %(#.(/; %<'5k
- -[( M%1%C#3@%<'( #.( '2%( 5#.'F/"%( .#"( 8A5'"A>7'%8( "%3"%5%<'/'A#<( #.( &#@@#<( /<8( 53%&A.A&P<#FC%8; %(#.(3/"'A&7C/"(/; %<'5(/<8(A<.#"@/'A#<(5%87"A'9(595'%@(/5(/(F2#C%K
- O[(D#8%CA<; (/<8(3)"'A/C(A@3C%@%<'/'A#<(#.('2%(1/5%(6'789(#.('2%(/; %<'=>/5%8(A<.#''@/'A#<5%87"A'9(595'%@4

: &&#"8A<; ('#('2%(8%1%#3%8('%&2<#C#; 9(5733#"'%8(>9(D: 6Mb('2%(@7C'A=/; %<'(595'%@(@#8%CA<; /<8(A@3C%@%<'/'A#<(#.(/; %<'=>/5%8(1\_66(5733#5%('#(5#C1%('F#(2A; 2(C%1%C'/5P5B(Z\*[(8%1%C#3@%<'#.('2% System Kernel #.('2%(1\_66E(/<8(Z-[(Cloning of the software agents(&#@3"A5A<; (1\_66(/<88%'/&2@%<'(#.('2%(; %<%''/'%8(@7C'A=/; %<'(595'%@(."#@(695'%@(b%''<%K4

O2%(.A</C(32/5%(#.('2%(! "#\$%&'("%5%/"&2(52#7C8(>%('2%(8%1%\#3@%<'(/<8(A@3C%@%<'/'A#<(#.('2% 1/5%(6'789E(F2A&2(2/1%('#(A<'%; "/'%('2%(>/5A&('2%#"%'A&/C(/<8(/"&2A'%&'7"/C(A8%/5(8%1%\\#3%8(/'3"%1A#75(32/5%5(#.('2%("%5%/"&24(O2%(>%\\#F(@/'%"A/C5(\#7'CA<\%5('2%(8%1\\\\\#3\\8(1/5\\)6'7894

#### 1.3. High-level Architecture of the Case Study

V%#F(A<'\2%(&/5%(5'789(F%(&#<5A8%"('2%(&#@37'%"(<%'F#"P(&#<5A5'A<; (#.(.#7"(2#5'5(ZYA; 4\*4\*[4 O2A5(<%'F#"P(&/<(&#@3"A5%(5%1%"/C(5%; @%<'5(#.(/(`:\_\_(/<8('2%(A<37'('"/..A&(&/<(>%(>#'2(A<5A8%(/<8 #7'5A8%(`:\_\_('"/..A&4( R/&2( 2#5'( A5( 3"#'%&'%8( >9( /< ( A<'%; "/'%8( /; %<'=>/5%8( 5%&7"A'9( 595'%@ &#@3"A5A<; (5%1%"/C(53%&A/CA]%8( agent-demons( I: M[( @#<A'#"A<; ('2%(A<37'('"/..A&( /<8( %S'"/&'A<; 3"%8%.A<%8(5A; <A.A&/<'(%1%<'5(/<8(3/''%"<5(."#@('2%(5%U7%<&%(#.(IP=3/&P%'5E(access control(/; %<'E authentication and identification(/; %<'(/<8( P<#FC%8; %=>/5%8( intrusion detection(/; %<'54( O2%5% /; %<'5(A<'%"/&'(87"A<; ('2%1"(#3%"/'A#<(1A/(@%55/; %5(%S&2/<; %('275(5733#"'A<; (&##3%"/'A1%(@7C'A=C%15%(B/\*)'/(3"#&%55A<; 4

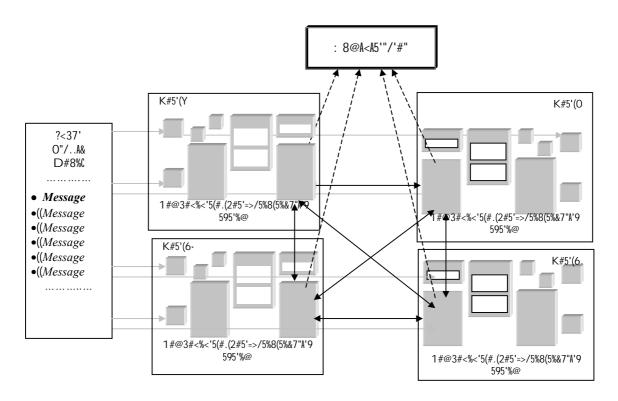


Fig.1.14(h33%"(C%I %C(/"&2A'%&'7"%(#.('2%(&/5%(5'789

Performative: tell

Sender: Input Traffic Model
Receiver: AD-E (host S1)

Content: <content>
Ontology: ISS
Language: XML

 $O2\%(\&\#<'\%<'(\#.('2A5(@\%55/;\%(A<(e\,D^{(C/<;7/;\%(A5(3"\%5\%<'\%8(A<('2\%(733\%'=C\%.'(\&\#''<\%''(FA<8\#FA<(YA;4^*-4$ 

O2%"%( A5( <#( 3/"'A&7C/"( 733%"=C%I %C( @/</; A<; ( 5#.'F/"%( '2/'( &##"8A</'%5( 2#5'=>/5%8( /; %<' #3%"'/'A#<( /5( F%CC( /5(A<'"75A#<( 8%'%&'A#<( /; %<'5( #.( 8A..%"%<'( 2#5'5( #.( '2%( &#@37'%"( <%'F#"P4( O2%@7C'A=/; %<'(5%&7"A'9(595'%@(#3%"/'%5(A<('2%(8A5'"A>7'%8(F/94

O2%(\?<37'(O"/..A&(@#8%C\(5#.'F/"%(Z5%%(YA; 4\*4\*[(5A@7C/'%5('2%(A<37'('"/..A&4(O2%(#7'37'(#. '2A5(3"#; "/@(A5(/("/<8#@(5%U7%<&%(#.(IP-3/&P%'5('2/'(A5(/(@AS'7"%(#.(<#"@/C(/<8(/><#"@/C(5'"%/@5#.(%1%<'54

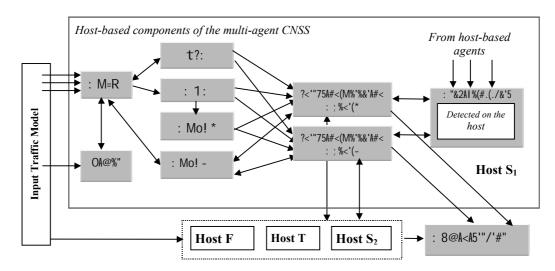
O2%(/><#"@/C(5"%/@(#.(%1%<'5(A5(5A@7C/'A<; (/''/&P5(/; /A<5'('2%(&#@37'%"(<%'F#"P4(Y#"(%1%"92#5'('2%(5#.'F/"%(&#@3#<%<'(\?<37'(O"/..A&(D#8%C\(A5(; %<%"/'A<; ('2%(A<37'('"/..A&(A<('2%(.#"@(#.5%U7%<&%(#.(#7'37'5(#.('2%(tcpdump(3"#; "/@4(Y#"(%S/@3C%E(5A@7C/'%8(SYN-port scanning(A<37'5%U7%<&%(&#7C8(>%(/5(.#CC#F5B

: <(/'''\&P(\'('2\%(\/33CA&\'A#<\(C\%1\%CE(\/("\%@#'\%(\/''\&P(\%S\&7'A<;\(\&#@@\/<85(\#.\('2\%(\#3\\"'\A#<\595'\\@(\#"\(\%S\&87'A<;\(\/33CA&\'A#<\(595'\\@(\&\CC5(\/"\%(53\&&A.A\\8(\A<\'\\@5(\#.\(\/(5\\U7\\<&\\%\\#.\('2\\(tcpdump\)@\%55\);\%5\\%S3\/<8\\8(\>9\'2\\\\\\Data Field\\(F2\&2\(53\\&A.A\\5\))'\2\\(\@\\%55\);\%(5\\@\/<\'\A\\5\)

O2% (5#.'F/''% (\$#8% (A5(F"A''%<(#<('2%(>/5A5(#.(WA57/C(1ccE(d:W:(-(/<8(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/<8(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/<8(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/<8(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/<8(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/48(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/48(eD`(5#.'F/''% 8%1%C#3@%<'(PA'54)A5(#.(WA57/C(1ccE(d:W:(-(/48(eD`(5#.'F)A5(#.(WA57/C(1ccE(d:W:(-(/48(eD`(1ccE(d:W:(-(/48(eD`(1ccE(d:W))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A5(A5(eD`(1ccE(d))A

#### 1.4. More Detailed Architecture of Host-based Security Components

`%'(75(&#<5A8%"(/"&2A'%&'7"%(#.('2%(2#5'=>/5%8(3/"'(#.('2%(@7C'A=/;%<'(1\_\_66(ZYA;4\*4-[4(O2% >/5A&( 2#5'=>/5%8( &#@3#<%<'5( /"%( /;%<'=8%@#<5( 8%<#'%8( >9( tM=RE( : M=!\*E( : M=!-E( : ?: Z/7'2%<'A&/'A#<(/<8(A8%<'A.A&/'A#<(/;%<'[(/<8(: 1 : (Z/&&%55(&#<'"#C(/;%<'[(/<8('F#(P<#FC%8;%=>/5%8 /;%<'5(=(?M: \*(/<8(?M: -(Z?<'"75A#<(M%'%&'A#<(: ;%<'5[4



**Fig.1.2.**(: "&2A'%&'7"%(/<8(A<'%"/&'A#<(#.(2#5'=>/5%8(3/"'(#.('2%(@7C'A=/; %<'(1\_66

 $Agent-demon\ AD-E(\ A5(\ ``\%53\#<5A>C%(\ .\#''(\ '2\%(\ A<37'(\ ''')...\ A&(\ 3"\%3"\#8\%55A<;\ 4(\ ?'(\ @\#<A'\#"5(\ ''')...\ A&(\ A@A<;\ (\ /'(\ \%5''')&'A\#<(\ \#.(\ 5\%U7\%<8\%5(\ \#.(\ 5\#=\&/CC\%8(\ `\%1\%<'5\(\ '2/'(\ /'\%(\ 5\%@/<'A&/CC9(\ @\%/<A<;\ .7C(\ /<8)).$ 

Agent-8%@#<5(.#"(A8%<'A.A&/'A#<(/<8(/7'2%<'A&/'A#<(Z: ?: [(/<8(/&&%55(&#<'"#C(Z: 1: [(3%".#"@'2%A"(&#<1%<'A#</Ci.7<&'A#</Ci.7<&'A#</Ci.7<&'A#</Ci.7<&'A%5E("%&#"8('2%("%57C'5(A<'#('2%A"(8/'/<5%5(/<8(5%<8(@%55/; %5('#('2%Intrusion detection agent(A.(/(5753A&A#75(>%2/1A#"(#"(/''%@3'(#.(/</''/&P(2/5(>%%<(8%'%&'%84(: ?: /; %<'(A5("%53#<5A>C%(.#"(A8%<'A.A&/'A#<(#.('2%(@%55/; %5(5#7"&%(/88"%55(/<8(.#"(&#<.A"@/'A#<(#.('2%A"/7'2%<'A&A'94(: 1: (/; %<'(@/</; %5(75%"5i(/&&%55(A<(&#<.#"@A'9(FA'2('2%A"(/&&%55("A; 2'5(/<8(C/>%C5(#. &#<.A8%<'A/CA'9(#.('2%(3"#'%&'%8(#>\$%&'54(O2%5%(/; %<'5(8%'%&'(/''%@3'5(#.('2%(7</7'2#"A]%8(/&&%55('#'2%(2#5'(A<.#"@/'A#<("%5#7"&%5E(8%'%&'5(A<'%""73'(%1%<'5(\Authentication failure" /<8 "Unauthorized access"E('"/3('2%5%(%1%<'5(/<8(A<.#"@(/>#"@(/>#7'(57&2(%1%<'5('2%(?M): -4

Agent-demons( AD-P1 /<8 AD-P2( /"%( "%53#<5A>C%( .#"( %S'"/&'A#<( #.( '2%( 3"%8%.A<%8 Z\A<'%"%5'A<; \[( \@%/<A<; .7C\[( 3/''%"<5( #.( \%1 %<'5\( /<8( @/PA<; ( 8%&A5A#<5( Z>#'2( 3"%CA@A</"9( /<8 .A</C[( />#7'( /''\%@3'5( #.( /'''/\&P5( #"( 5753A&A#75( >%2/1 A#"( #.( 75%"5( &#<<%&'\%8( FA'2( '2%( 2#5'4(?<( '2% A@3C\%@%<'\%8(&/5\%(5'789('2\%/', %<'(: M=! \*(A5("\%53#<5A>C\%( .#"(\%5'"/\&'A<; (3/''\%'<5('2/'(&#""\%53#<85('#

'2%(5753A&A#75(>%2/1A#"(#"(/''%@3'5(#.(/''/&P5(CAP%(port scanning(Z#<(/33CA&/'A#<(C%1%[E(finger search(/<8(buffer overflow4(O2%(/; %<'(: M=! -(A5(A<'%<8%8('#(%S'"/&'(3/''%"<5(&#""%53#<8A<; ('#('2% denial of service(/'''/&PE(syn-flood(/'''/&P(/<8(port scanning(Z#<(<%'F#"P(/<8('"/<53#"'(C%1%[4

O2%(5753A&A#75(%1%<'5(/<8(3/''%''<5(8%'%&'%8(>9(/; %<'=8%@#<5(/''%(5%<'('#(?M: \*(/<8a#"'(?M: -/<8(5#@%(#.('2%@(/"%("%&#"8%8(A<'#('2%(p: "&2A1%(#.(../&'5q4(?M: \*(/<8(?M: -(75%('2%5%(../&'5('#3%".#"@(P<#FC%8; %=>/5%8(3"#&%55A<; (#.('2%(A<.#"@/'A#<("%&%A1%8(."#@(/; %<'=8%@#<54

Database "Archive of facts" (5"#"%5(8/'/(>#7'(3/5'(72/1%(>%<(#1%"[(&#<<%&'A#<54(O2%5%(8/'/"%(75%8(>9(?M: \*(/<8(?M: -(.#"(8%'%&'A#<(8A5'"A>7'%8(@7C'A=32/5%(/''/&P5(F2%<(A'(A5(<%&%55/"9('#3"#&%55(&7""%<'(/<8(3/5'(&#<<%&'A#<5('#("%&#1%"('2%(5&%</"A#(#.(/''/&P(/<8(@/'&2('#(/'(P<#F<5&%</"A#4(\_#'A&%('2/'(8%'%&'A#<(#.(/(5&%</"A#(#.('2%(/''/&P(A5(3#55A>C%(#<C9(/5(/("%57C'(#.(&##3%"/'A#<#.(5%1%"/C(/;%<'54

- Z\*[(ReconnaissanceE('2/'(&#""%53#<85('#('2%(&#CC%&'A<;(;%<%"/C(8/'/(/>#7'(&#@37'%"(<%'F#"P &#<.A;7"/'A#<E( A8%<'A.A&/'A#<( #.( '2%( 2#5'5( &#@3"A5A<;( <%'F#"PE( /&&%55A>C%( 5%"IA&%5E A8%<'A.A&/'A#<(#.('2%(#3%"/'A#<(595'%@E(/33CA&/'A#<5(%'&4
- IO[(Privileges escalating('2/'(A5(/(@/C%./&'#"i5(/''%@3'('#(;%'(%S3/<8%8(/&&%55("A; 2'54
- ZG[(Deepening penetration(#<(2#5'('2/'(A5(%S3/<8A<;(@/C%./&'#"i5(/&&%55('#(.AC%5E(.#C8%"5(/<83"#;"/@5(#.('2%(2#5'
- ZN[(Deepening penetration through net('2/'(A5(%S3/<8A<; (@/C%./&'#"i5(/&&%55('#(.AC%5E(.#C8%"5(/<8 3"#; "/@5(#.('2%(#'2%"(2#5'5(#.('2%(&#@37'%"(<%'F#"P(7<8%"(/''/&P4
- V#'2(?M: \*(/<8(?M: -(2/1%(5A@AC/"(/"&2A'%&'7"%(/<8('2%(#<C9(8A5'A<&'A#<(A5('2%(&#<'%<'(#.('2%A" P<#FC%8; %(>/5%5(Zb)V[(&#@3"A5A<; ('2%("7C%5E(75%8('#(@/P%(8%&A5A#<5(#<('2%(>/5A5(#<('2%(>/5A5(#<('2%(>/5A5(#</\*\*).454)
- \_#'A&%('2/'(%/&2(/; %<'(#.('2)(2)#5'(2/1)%(/(C)#&/C(8/')>/5%(F2A&2(A)5(/&&%55A)>C)(.#"(%1\%"9(2#5'= >/5\%8(/; %<'4(?<(./&'E(/C)#; %'2\%"('2\%9(.#"@(8A5'"A>7'\%8(8/')>/5\%(#.('2\%(2\#5'=>/5\%8(5\%87"A'9(595'\%@4
- O2%( .7<&'A#</CA'9( #.( '2%( 5#.'F/"%( &#@3#<%<'( \: 8@A<A5'"/'#"\( A5( '#( 3"#1A8%( /( 75%"( >9 4.0M

#### 1.5. Agent Architecture

: CC( 2#5'=>/5%8( 5%&7"A'9( /; %<'5( /'%( A@3C%@%<'%8( #<( '2%( >/5A5( #.( /( 5'/<8/"8( /"&2A'%&'7"% 8%3A&'%8(A<(YA; 4\*404(O2%(8A5'A<&'A#<5(/'%(#<C9(A<(P<#FC%8; %(>/5%5(&#<'%<'54(O2%(/; %<'(8%@#<5i(b)V &#<'/A<(#<C9(5A@3C%("7C%5(.#"(%S'"/&'A<; (%1%<'5(#"(3/''%"<5(F2%"%/5(A<'"75A#<(8%'%&'A#<(/; %<'5(2/1%@#"%(&#@3"%2%<5A1%(b)V54(` %'(75(#7'CA<%('2%(; %<%"A&(/"&2A'%&'7"%(#.(/; %<'54

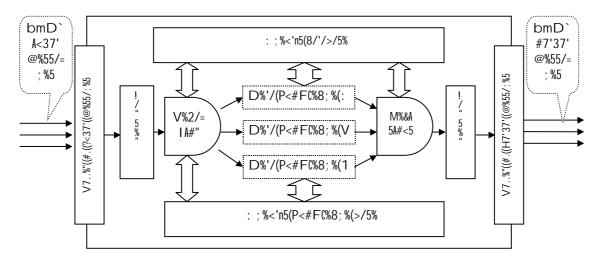
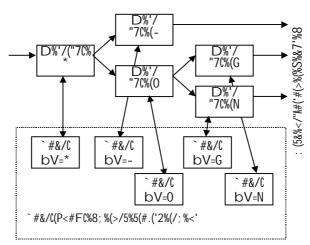


Fig.1.3.(6%&7"A'9(/; %<'(/"&2A'%&'7"%

 $Buffer\ of\ input\ messages (A5(75\%8(.#"(A<37'(@\%55/; \%5(5'#"A<; E(#"8\%"A<; (/<8(59<\&2"\#<A]/'A\#<(#. '2\A"(3"#&\%55A<; 4(`\%'(75("\@A<8('2/'(/\C(@\%55/; \%5(/"\%("\%3"\%5\<'\\8(A<(\bmD`(.#"@/'E(F2#5\%\&#<'\%<' A5(53\&A.A\\8(A<('\%''@5(#.(\ed)^`(\c/<; 7/; \%4$ 

Parser(3%".#"@5(59<'/&'A&/C(/</C95A5(/<8(A<'%"3"%'/'A#<(#.(@%55/;%54(O2%(5A@AC/"(3"#;"/@3%".#"@5(;%<%"/'A#<(/<8(5%<8A<;(#.(#7'37'(@%55/;%54(O#(A<'%"3"%'('2%(A<37'(/<8('#(;%<%"/'%('2%#7'37'(@%55/;%54(O#(A<'\*"3"%'('2%(A<37'(/<8('#(;%<%"/'%('2%#7'37'(@%55/;%(!/"5%"(2/5('#(/&&%55('#('2%(8#@/A<(#<'#C#;9(#.('2%(/;%<'E(F2A&2(A5(/(3/"'(#.('2% Agent's database4



**Fig.1.4.**(: <(%\$/@30%(#.('2%(/; %<'i5 P<#F0%8; %(>/5%(/"&2A'%&'7"%

O2%(A<'%"3"%'%8( @%55/; %( &#<'%<'( A5( '2%( A<37'( #.( '2%( &#@3#<%<'( Behaviort( F2A&2( 2/5( '#8%'%"@A<%( '2%( F/9( #.( A<37'( @%55/; %( 3"#&%55A<; 4( ?'( ; %<%"/'%5( /( 5&%</"A#( #<( '2%( >/5A5( #.( '"%%5'"78'7"%8( Meta-knowledge baset( F2A&2( A.( <%&%55/"9( 8%"A1%5( /( 8%&A5A#<( 1A/( /&&%55A<; ( Local

: ("%53#<5A>ACA'9(#.('2%(Decision(&#@3#<%<'(A5('#(8%'%@A<%(/("%/&'A#<(#.('2%(/; %<'(A<('2%'%57C'(#.('2%(5&%</"A#(%S%&7'A#<4(O2%'%(/"%('\*F#(3#55A>C%("%/&'A#<5B(Z\*[(; %<%''/A#<(#.('<(#7'37'@%55/; %Z5[(/<8(Z-[("%&#"8A<; ('2%("%57C'(#.('2%(5&%</"A#(%S%&7'A#<(A<('2%(/; %<'i5(8/')>/5%4(O2%(C/''%'')''5(/("70%(&#""%53#<85('#(/<(A<'%''@%8A/'%(8%&A5A#<(@/8%(>9('2%(/; %<'('2/'(A5(5733#5%8('#(>%(75%8 C/'%''4

#### 1.6. Input Traffic Model: Scenarios of Combined Distributed Attack

O2%( @#8%( #.( A<37'( '"/..A&( &#""%53#<85( '#( /( <7@>%"( #.( 5&%</"A#5( '2/'( /"%( 75%8( >9( /@/C%../&'#"4(: (5&%</"A#(&/<(&#<5A5'(#.(/(57>5%'(#.(/''/&P5('2/'(/'"%

\*4(6&/<<A<; (#.(/(2#5'(3#"'5(/A@A<; (/'(8%'%"@A</'A#<(#.('2%(/&'AT%(3#"'5(/<8(/T/AC/>C%(5%"TA&%5(Z.'3E"'%C<%'E(%'&4[4

-4(: ''%@3'5(#.(&#<<%&'A#<('#(/(2#5'(#<(/(5%"1A&%(CAP%(.'3E('%C<%'E(%'&4(75A<; (8A..%"%<'(5#7"&%(?! = /88"%55%5(/<8(; 7%55A<; (3/55F#"84

O4(1#@>A<%8(53##.A<; (/''/&P4(O2A5(/''/&P(/; /A<5'(/(&#@37'%"(<%'F#"P(&#@3"A5%5(5%1%"/C(5"%354 YA"5'E('2%(@/C%./&'#"(%5'/>CA52%5(/(IP=&#<<%&'A#<(#<(/(2#5'(6(/<8(; %'5(/<(A<A'A/C(<7@>%"(#.('2% &#<<%&'A#<4(\_\_%S'E(2%("%/CA]%5(denial of service(/''/&P(/; /A<5'(/('"75'%8(2#5'(O(/<8(87"A<; (O 2/<; =73(2%(%5'/>CA52%5(/(IP=&#<<%&'A#<(#<('2%(2#5'(6(#<(>\*2/C.(#.('2%('"75'%8(2#5'(O4(O2%(2#5'6(5%<85('#('2%(2#5'(O('2%(2#5'(O2%(2#5'(O2%(2#5'(O('3#55%55%54')

G4(: ''%@3'5('#(;%'(7</7'2#"A]%8(/&&%55('#(.AC%5(#.(/(2#5'4

N4(V7..%"(#1%".C#F(/''/&P(#<(/(2#5'(/A@A<; (/'(>##5'A<; (#.('2%(/&&%55("A; 2'54

?<(./&'E(\?<37'(O"/..A&(D#8%C\(A5(/(8/'/>/5%E(F2A&2(A5(5'#"A<; ('2%(5%U7%<&%(#.(IP=3/&P%'5('2/' A5( 5733#5%8( '#( >%( 75%8( /5( '2%( A<37'( #.( /( 2#5'4( : ( 3/"'A&7C/"( 3"#; "/@( '2/'( A5( 8%1%E#3%8( /<8 A@3C%@%<'%8(/5(/(5#.'F/"%('##C(5A@7C/'A<; (<#"@/C(/<8(/><#"@/C(A<37'('"/..A&(; %<%"/'%5('2%(&#<'%<' #.('2A5(8/'/>/5%4(h33%"=C%1%C(75%"(A<'\%"./&%(#.('2A5('##C(A5(8%3A&'%8(A<(YA; 4\*4N4

`%'(75(%S3C/A<(A<(>"A%.(2#F(/(@#8%C(#.(/(&#@>A<%8(/''/&P(A5(&#@3#5%84(O2%(.A"5'(5'%3(A5('#8%.A<%('2%(5%U7%<&%(#.(\2%5%5)#.(\5%3(A5('#8%.A<%('2%(5%U7%<&%(#.(\2/@%5(#.(\5A<);C%32/5%\(/''/&P54(O2A5(5%U7%<&%(8%'%"@A<%5('2%(2A;2(C%1%C(53%&A.A&/'A#<(#.('2%(/''/&P('#(>%(5A@7C/'%84O2%(<%S'(5'%3(#.(/''/&P(&#@3#5A'A#<(A5('#(@/3(/(5%U7%<&%(#.(IP=@%55/;%5('#(%/&2(</@%(#.(\5A<);C%=32/5%(/'''/&P\(8%'%"@A<%8(A<('2%(/.#"%@%<'A#<%8(5&%</"A#4(IP=@%55/;%5(2/1%('#(>%(%8A'%8('#(/55A;<'2%''A@%(#.(%&2(\5A<);C%=32/5%\(/'''/&P(>%;A<)('''A\$('2%'(/''''A>7'%54(O2%(5%&#<8(5'%3(#.(%8A'A#<(A5('#

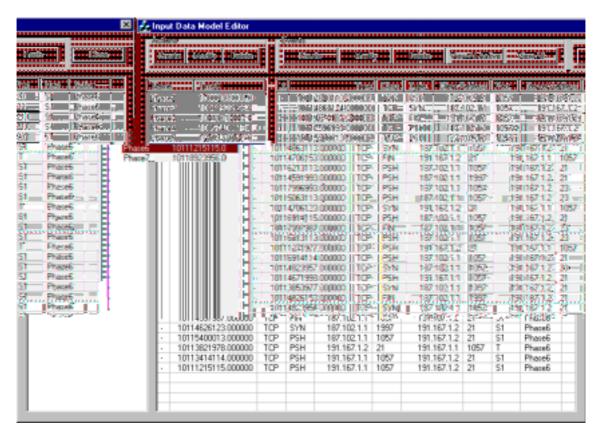
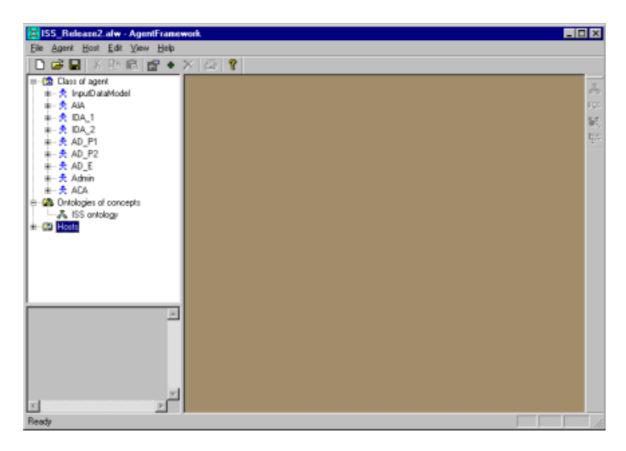


Fig.1.5.(h5%"(A<'%"./&%(R8A'#"(#.('2%(A<37'('"/..A&(@#8%('#(>%(5A@7C/'%8



**Fig.1.6.**(`A5'(#.(/; %<'(&C/55%5(#.('2%(1/5%(6'789

#### 1.7. Technology of the Case Study Development and Implementation

 $O2\% (1/5\% 6'789(\#.('2\% @7C'A=/; \%<'(<\%'F\#"P(5\%\&7"A'9(595'\%@(8\%5\&"A>\%8(A<('2\% 3"\%1A\#755\%\&'A\#<5(F/5(8\%1\%2\#3\%8(/<8(A@3C\%@%<'\%8(/5(/<(/33CA&/'A\#<(\#.('2\%(5\#.'F/"\%('\#\#C(\D7C'A=/; \%<'695'\%@(M\%1\%2\30\%2'(bA'\(ZD:6Mb[4(:7'2\#"5(\#.('2A5(J\33\#"'(8\%1\%2\30\%8(5\#.'F/"\%('\##C 8\%5\%"A>\%8(A<('2\%3'\%1A\#75(!\#\$\%\'(J\33\#"'(FYA<J\33))*q4$ 

First Stage. Development of the System Kernel.

#### 1. Designation of the agent classes of the Case Study

#### 2. Ontology creation

6AS(&#<&%3'5(/"%(A<&C78%8(A<'#('#3=C%1%C(8#@/A<(#<'#C#; 9(75%%(YA; 4\*4Q[4(: <(%S/@3C%(#.(/(>/5A& &#<&%3'(#.('2%(1/5%(6'789(&/CC%8(\17""%<'(1#<<%&\A#<\\(A5(8%5&\"A>%8(A<('2%(: 33%<8AS4

#### 3. Designation scenarios to each class of the agents

?<('2%('/>c%(\*4\*(>%C#F('2%(CA5'(#.('2%(>%2/I A#"/C(5&%</"A#5(/&&%55A>C%(.#"(%/&2(/; %<'(A5(; A1 %<4(O2% 5&%</"A#5(/"%(8%5&"A>%8(A<('%"@5(#.(#3%"/'A#<5('2/'(/"%(&/""A%8(#7'(>9(/; %<'4(O2%(&#@3C%'%(CA5'(#. 57&2(5&%</"A#5(A5(; A1 %<(A<(: 33%<8AS(: \*4

O/>C%(\*4\*4(! /"'A/C(CA5'(#.(/; %<'5i(>%2/I A#"/C(5&%</"A#5

: ; %<'(&C/55%5	68% "A#5(#.( %2/1A#"
: M=R	• ! "%3"#&%55A<; (#.(tcp=@%55/; %5(#.(A<37'('"/A&
	• 12%&PA<; (\\life time\(#.('2%(&#<<%&'A#<5(>%C#<; A<; ('#('2%(&7""%<'(&#<<%&'A#<5</td></tr><tr><td></td><td>• ! "#&%55A<; (#.(8%'%&'%8(./&'5(A<('2%(&7""%<'(&#<<%&'A#<(8%5&"A3'A#<</td></tr><tr><td>: M=! *</td><td>• 12%&PA<; (#.('2%(&7""%<'(<7@>%"(#.('2%(#&&7""%<&%5(#.(\5753A&A#75\(>%2/IA#" /<8(8%'%&'A#<(#.('2%(Port scanning /<8 Syn Flood /''/&P5</td></tr><tr><td>: M=! -</td><td>M%'%&'A#<(#.(Port scanning{(Buffer overflow(/<8(Finger search))</td></tr><tr><td>: ?:</td><td>• : 7'2%<'A&/'A#<(/<8(A8%<'A.A&/'A#<</td></tr><tr><td>: 1:</td><td>• : &&%55(&#<'"#C</td></tr><tr><td>?M: =*</td><td>M%'%&'A#<(#.('2%(Combined spoofing attack))</td></tr><tr><td>?M: =-</td><td>• b<#F0%8; %=>/5%8(3"#&%55A<; (#.('2%(8%'%&'%8(./&'5E(3/''%"<5E(%'&4</td></tr><tr><td>: 8@A<A5'"/'#"</td><td>• 6'#"A<; (/<8(1A57/CA]/'A#<(#.('2%(8/'/(/>#7'(&#<<%&'A#<5(/<8(/>#7'(8%'%&'%8 /''/&P5</td></tr></tbody></table>

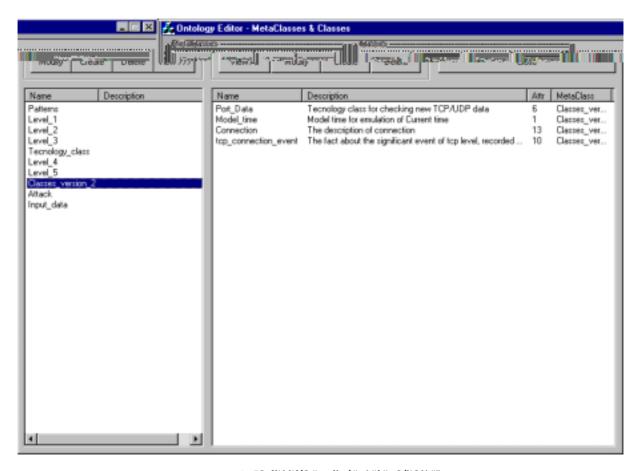


Fig.1.7.(O#3=C%I %C(8#@/A<(#<'#C#; 9(%8A'#"

Second Stage: Cloning of the software security agents

O#(>%; A<('2%(&C#<A<; (3"#&%87"%(A'(A5("%U7A"%8('#(8%'%"@A<%('2%(</@%5(#.(2#5'5(A<(F2A&2('2%(5%&7"A'9595'%@(A5(5733#5%8('#(>%(5A'7/'%8(ZYA; 4\*4, [4(?'(52#7C8(>%(<#'A&%8('2/'(A.(F%(2/I%('#(@/P%(8A5'A<&'A#<>\*\'F%%<('2%(</@%5(#.(2#5'5(comprising the Case Study(/<8('2%(</@%5(#.(computers in which the Case Study is supposed to install4(O2%(1/5%(6'789(A5('2%(@#8%C(#.(/(5%&7"A'9(595'%@(/<8(A'(&/<&#@3"A5%(N 2#5'5('#(>%(3"#'%&'%84(Y#"(%S/@3C%E(A<(#7"(&/5%(Z5%%(YA; 4\*4\*[(A'(&#<5A5'5(#.(.#7"(2#5'5<//@%8(YE(OE(6\*E(/<8(6.4

`%'(75(8%<#'%('2%(5%'(#.(/; %<'5('2/'(/&&#"8A<; ('#('2%(8%1 %#3%8(1/5%(6'789(#.('2%(@7C'A=/; %<'5%&7"A'9(595'%@(/"&2A'%&'7"%(@75'(>%(5A'7/'%8(#<(/(2#5'(j (>9(AZj [4(?<(#7"(1/5%(6'789('2%(%<'A"%(5%'#.(/; %<'5('#,0)%(6'789('2%(%<'A"%(5%'#.(/)2#5'(j (>9(AZj [4(?<(#7"(1/5%(6'789('2%(%<'A"%(5%'#.(/; %<'5('#,0)%(6'789('2%(%<'A"%(5%'#.()2%(0)AZ))\_{4(?)})\_{4(?)}\_{

-4(?<5'/CC/'A#<('2%(5#.'F/"%(*Agent cloning* A<(%/&2(#.('2%(/>#1%(&#@37'%"54(O2A5(5#.'F/"%(3"#1A8%5.#"("%@#'%(/&&%55(."#@('2%(&#@37'%"(A<(F2A&2('2%(695'%@(b%"<%(A5(A<5'/CC%8('#('2%(.AC%5(#.('2%(2#5'A<(F2A&2('2%(695'%@(b%"<%(A5(A<5'/CC%8('#('2%(.AC%5(#.('2%(2#5'A<(F2A&2('2%(695'%@(b%"<%(A5(A<5'/CC%8('#')'%(A<5'/CC%84(?<(/88A'A#<E(A<(%/&2(57&2(&#@37'%"(5%1'%"(/5#.'F/"%(/<8('2%(CA>"/"9(#.(; %<%"A&3"#&%87"%5('#(A@3C%@%<'(5&%</"A#5(#.(/; %<'i5(>%2/1A#"4

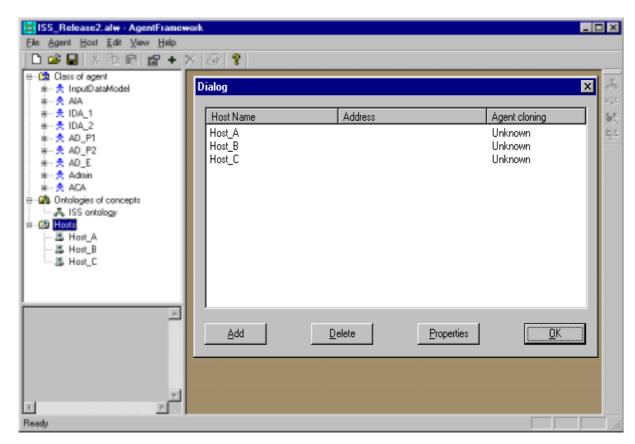


Fig.1.8.((O2%(2#5'(</@A<; (/<8("%/C(&#@37'%"(8%5A; </'A#<(A<'%"./&%

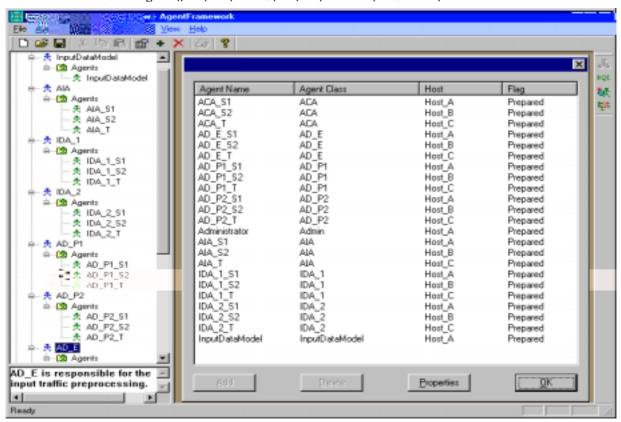
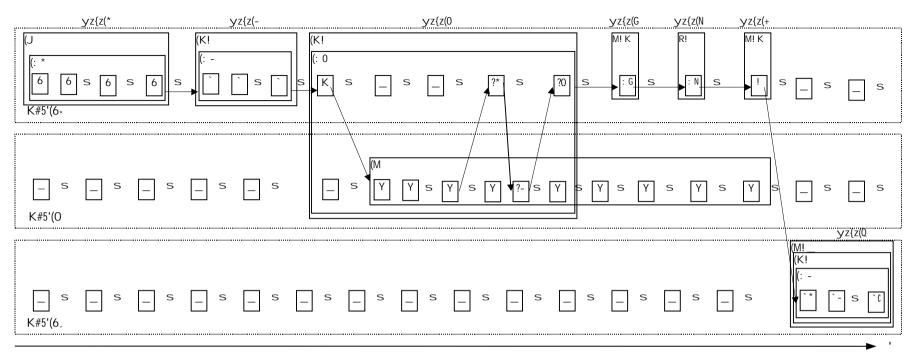


Fig.1.9.(63%&A.A&/'A#<(#.('2%(/"&2A'%&'7"%(#.('2%(@7C'A=/; %<'(5%&7"A'9(595'%@(('#(>%(&C#<%8

## 1.8. Case-study Simulation: An Example of a Distributed Attack and Multi-agent System Performance

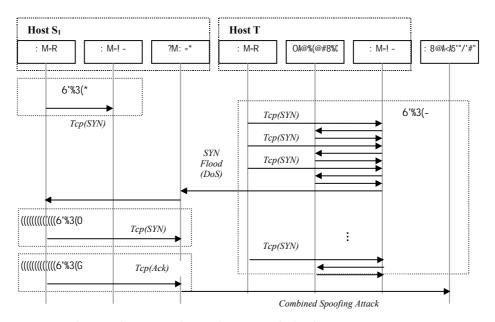
 $6A@7C/'A\#<(F/5(\#";/<A]\%8(\#<('2\%)</5A5(\#.('2\%(.\#FA<;;(5'''/\%;9(\#.(/'''/\&P\%'4(?'(A5(5733\#5\%8''2/'('2\%(/'''/\&P(5\#7''8\%(A5('2\%(@/C\%./\&'#"i5(2\#5'(<math>\oplus$ (/<8(/'''/\&P(A5('''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5(''')/\&P(A5('')/)/P(A5('')/\&P(A5('')/\&P(A5('')/)/P(A5(



Particular events: (6Es E6(0(6j \_\_(0(3/&P%'5(Z.#"(6j \_\_=5&/<<A<; [E(`Es E`(0(C#; A<(N1%<'5E(&#""%53#<8A<; ('#('2%(/''%@3'5(#.(/(3/55F#"8(; 7%55A<; E(K(0(%1%<'(K/<852/P%("%&#<</A55/<&%EYES EY(0(6j \_\_=3/&P%'5(Z.#"(6j \_\_(YC##8[E(?\*(0(%1%<'5(&#""%53#<8A<; ('#('2%(A®A')'\A#<(#.('2%(A°5'(32/5%(#.(2/<852/PA<; E(?-(0(&#<.A"@/'A#<(3/&P%'(."#@('2%(2#5'(6-E(?0(0(%1%<'(&#""%53#<8A<; '#('2%(A@A')'\A#<(#.('2%(2A)'\A\*\*(#.('2%(2A)'\A\*\*(#.('2%(2A)'\A\*\*(#.('2%(A)'\A\*\*(#.('2%(A)'\A\*\*(#.('2%(A)'\A\*\*(#.('2%(A)'\A\*\*(A)'\A\*\*(#.('2%(A)'\A\*\*(A)'

X%<\"/CA]\%8(/''\&P(5&\%</"A\*5\(\_)o\"\&#(\K!\)o\2\\$5'(3\%<\"'\\^\\#\<\[(\K!\)o\2\\$5'(3\%\\"\'\\\#\<\[(\K!\)o\2\\$5'\[(\K!\)o\2\

Fig.1.10.(OA@%=#"8%"%8(8A/; "/@(#.('2%(3/"'A&7C/"(/''/&P(5&%</"A#



**Fig.1.11.**(: ; %<'5i(A<'%"/&'A#<(87"A<; (3"#&%55A<; (#.('2%(Combined Spoofing Attack

\*4(6&/<<A<; (#.('2%(3#"'5(#.('2%(2#5'(6\*4

-4(: ''%@3'5(#.(&#<<%&'\A#<(#<('2%(2#5'(6\*('#(/&&%55('#(A'5(.'3(#"('\%C<\%'(5\%"\A\&\%5(.75A<; (\/\"\A\#75(5\#7"\&\%\IP=/88\\%55\\%5(/<8(; 7\%55A<; (3/55\F\\\"\84(?'(\A\\$(5733\\\\5\\\\))''/\&P(\A\5(7<57\&\\\\\\5\\\)5.764

Z>[(D/C%./&'#"(e(&#@3"#@A5%5('2%(2#5'(O('2"#7; 2(8A"%&'%8(5'#"@(#.(@%55/; %5(Z6j \_=.C##8[('275 A@3C%@%<'A<; (\Denial of Service\(/''/&P4

Z8[(D/C%./&'#"(e(3#55%55A<; (>9('2%(&#<<%&'A#<(<7@>%"(?6\_5-("%3CA%5('#('2%(2#5'(6\*(A<5'%/8(#. '2%(2#5'(O('275(; %''A<; (5A<; C%(5A8%(&#<<%&'A#<(FA'2('2%(2#5'(6\*4(?<('2%(5A@7C/'A#<(3"#&%87"%(A' F/5(5733#5%8('2/'('2A5(/''/&P(F/5(57&&%55.7C4

+4(J%/8A<; ('2%(3/55F#"85(/<8(; %'`A<; (/&&%55('#('2%(.AC%5(#.('2%(2#5'(6\_(Z57&&%55.7C[4 Q4(1#<<%&'A#<('#('2%(.'3(#"('%C<%'(5%"I A&%(#.('2%(2#5'(6\_7Z57&&%55.7C]4

 %'(75(&#<5A8%''('2%(/; %<'i5(A<'%''/&'A#<5(87"A<; ('2%('2A"8(32/5%(#.('2%(/>#1%=8%5&"A>%8(/'''&PEF2A&2(&#'''%53#<85('#('2%(&#@>A<%8(53#.A<; (/'''&P(ACC75'''/'%8(A<(YA; 4\*4\*\*4))))))))))))

At the first step '2%(2#5'( $\ominus$ (%5'/>CA52%5('2%(tcp=&#<<%&'A#<(FA'2('2%(2#5'(6-4(?.('2%(&#<<%&'A#<FA'2('2%(2#5'(6-4(8)(%5'/>CA52%8('2%(tcp=@%55/; %(/55A; <%8(.C/; (6j \_\_(A5(5%<'(."#@('2%(/; %<'=8%@#<: M=R(#.('2%(2#5'(6-('#('2%(/; %<'(: M=! \*(Z5%%(YA; 4\*4\*\*[4?'(A5(A@3%"/'A1 %('2/'(57&2(/(@%55/; %('#(/(2#5'@75'(>%(.#C#F%8(>9('2%(<%S'(5A@AC/"(@%55/; %(/55A; <%8(.C/; (: 1  $\bigcirc$ 4(?.(57&2(/(@%55/; %(A5(<#'(>%A<; "%&%A1 %8(>9('2%(2#5'(6-('<8E('2%"%.#"%E(<#'("%&%A1 %8(>9(A'5(/; %<'(: M=! \*('2%<('2%(C/''%"(@/P%5\*&#<&C75A#<('2/''('2A5(tcp=&#<<%&'A#<(A5(\2/C.=#3%<\4

At the third step('2%(/; %<'=8%@#<(: M=R(#.('2%(2#5'( $6*(5\%<85('2\%(@\%55/; \%('#(A'5(A<'"75A#<8\%'%&'A#<(/'; %<'(?M: *('#(A<.#"@(A'(/>#7'("%&%1A<; ('2%(@%55/; %(/55A; <%8(.C/; (6j __(."#@('2%(2#5'(O4)))))))))))))))))))))))$ 

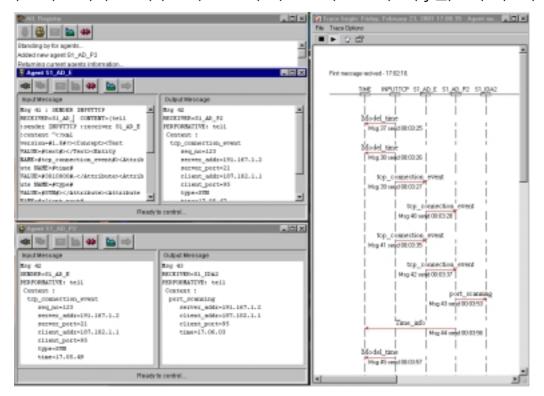


Fig.1.12.(! "A<'#7'(#.('2%(1A57/C(A<'%"./&%("%3"%5%<'A<; (/; %<'5i(A<'%"/&'A#<

: ; %<'(?M: \*(\$#A<5('2A5(./&'(FA'2('2%(./&'(A<.%"''%8(/'('2%('2A"8(5'%3(/<8(75A<; (A'5(P<#FC%8; %(>/5% &#<&C78%5('2/'(/(5753A&A#75(>%2/1A#"(A5('/PA<; (3C/&%(/'('2%(2#5'(A<37'4(?'(&C/55A.A%5('2A5(>%2/1A#"(/5 \possibly, unknown user imitates tcp-connection with the host  $S_1$  on behalf of the trusted host T\4

At the fourth step /; %<'=8%@#<(#.('2%(2#5'(6\*(5%<85('2%(@%55/;%('#('2%(A<'''75A#<(8%'%&'A#</; %<'(?M: \*(#.('2%(<math>5/@%(2#5'('2/'(A'(''%&%A1%8(@%55/;%(FA'2('2%(.C/; (: 1 b(."#@('2%(2#5'(O4(?M: \*A<.%"5('2/'('2%(&#@>A<%8(<math>63#.A<; (: ''/&P(A5('/PA<; (3C/&%)

?'(A5(F%CC(P<#F<('2/'(3/"'A&7C/"(./&'5(&/<(>%(%1A8%<&%5(#.(/(&#@>A<%8(8A5'"A>7'%8(/''/&P/;/A<5'(&#@37'%"(<%'F#"P4(O2%(5/@3C%(#.(57&2(/<(/''/&P(F/5(;A1%<(/>#1%(F2%<('2%('2%('2A"8(32/5%(#.'2%(@/C%./&'#"i5(5&%</"A#(#.(/''/&P(F/5(8%5&"A>%84(O2%(8%'%&'%8(32/5%5(#.(/(&#@>A<%8(/''/&P(/''%8%3A&'%8(A<('2%(C#F="A; 2'=2/<8(/"%/(#.('2%(8A/;"/@4(h5%"(A<'%"./&%(/CC#F5(@/"PA<; (Z>9(53%&A/C'''/&%"9[('2%(5%'(#.(./&'5(2/'(&#""%53#<85('#('2%(%1A8%<&%5(#.('2%(5/@%(/''/&P4



**Fig.1.13.**(h5%"(A<'%"./&%(.#"(/</C95A5(#.('2%("%57C'5(#.('2%(1/5%(6'789

#### 1.9. Conclusion of Chapter 1

: 5(/("7C%E(8%1%C#3@%<'(#.(/(8A5'"A>7'%8(/''/&P(&/<(/..%&'(5%1%"/C(2#5'5(&#@3"#@A5A<;(/ <7@>%"(#.('2%@(/'('2%(5/@%('A@%4(D/C%./&'#"5(&/<(&##3%"/'%(FA'2A<(/<(/''/&P(A<(57&2(/(F/9(A< F2A&2(/(3/"'A&7C/"(@/C%./&'#"(/''/&P5(/(5A<; C%(2#5'('275(@/5PA<; ('2%(&#@>A<%8(8A5'"A>7'%8(</'7"%(#. /<(/''/&P4(: (&%<'"/CA]%8(5%&7"A'9(595'%@(52#7C8(>%('##(\2%/19\('#(3"#&%55('2%(%<'A"%(>/'&2(#.(C/"; %= 5&/C%(8/'/(/>#7'(3"%2A5'#"9(/<8(&7""%<'(5'/'%5(#.(@/<9(&#<&7""%<'(&#<<%&'A#<5(#.(@/<9(2#5'54 6A@7C/'A#<(#.('2%(&/5%(5'789(8A53C/9%8(/(<7@>%"(#.(/81/<'/; %5(#.(@7C'A=/; %<'(/"&2A'%&'7"%(.#" 3"#'%&'A#<( #.( /( &#@37'%"( <%'F#"P( /; /A<5'( 8A5'"A>7'%8( /''/&P54( O2%( @#5'( 5A; <A.A&/<'( #<%( A5( / &/3/>ACA'9(#.(&#@3/"/'AI%C9(\CA; 2'\(&#@3#<%<'5(#.(/(@7C'A=/; %<'(5%&7"A'9(595'%@('#(&##3%"/'%4(: 3"%5%<'('2%(#<C9(F/9('#(8%'%&'(%..A&A%<'C9(/(8A5'"A>7'%8(/''/&P(/;/A<5'(/(&#@37'%"(<%'F#"P(A5(/ &##3%"/'A#<(#.(5%&7"A'9(5#.'F/"%(%<'A'A%5(Z/;%<'5[(8A5'"A>7'%8(#I%"('2%(2#5'5(#.('2%(<%'F#"P(/<8 FA'2A<(%/&2(2#5'\a'5%.4(: (5/@3C%(#.(75%.7C<%55(#.(//; %<'(&##3%"/'A#<(&/<(>%(5%%<(FA'2A<('2%(%S/@3C% 8%5&"A>%8( A<( '2%( 3"%1 A#75( 5%&'A#<4( ?<( 3/"'A&7C/"E( 8%'%&'A#<( #.( /( &#@>A<%8( 53##.A<; ( /''/&P( 2/5 >%&#@%(3#55A>C%(#<C9(87%('#(&##3%''/'A#<(#.(/; %<'5(: M=! -(/<8(?M: \*(#.('2%(2#5'(6\*(/<8(: M=! -(#...)) + ...)))))))) '2%( 2#5'( OE( A4%4( 87%( '#( /( &##3%"/'A#<( #.( 5#. 'F/"%( %<'A'A%5( 5A'7/'%8( #<( 8A..%"%<'( 2#5'54( H<%( @#"% %\$\@3C%(#.('2%(<%&%55A'9(\<8(75%.7C<%55(#.('2%(&##3%"\'A#<(#.(5%&7"A'9(\;%<'5(A5(#3%"\'A#<(#.('2% P<#FC%8: %=>/5%8(/; %<'(?M: -4(?'(A5(A<'%<8%8('#(&#CC%&'(A<.#"@/'A#<(/>#7'(5753A&A#75(>%2/IA#"(#. 75%"5(/'(@/<9(%<'"9(3#A<'5(#.('2%(</'AI'%(2#5'(/<8(%S&2/<; %(A<.#"@/'A#<(FA'2(5A@AC/"(/; %<'5(#.('2% #'2%"(2#5'5('#(@/P%(A<'%; "/C(8%&A5A#<(/>#7'(5'/'75(#.(&#<<%&'A#<54(?<('2%(8%1%C#3%8(&/5%(5'789('2A5 /;%<'(2/5(>%%<(3"#1A8%8(>9(/(&#@3/"/'A1%C9(3##"(P<#FC%8;%(>/5%('#(3C/9(/(5A;<A.A&/<'("#C%(A< A<'"75A#<(8%'%&'A#<4(?<('2%(.7"'2%"(8%1%C#3@%<'(#.('2%(&/5%(5'789(/(5A; <A.A&/<'(/&&%<'(FACC(>%(@/8% #<(%<"A&2@%<'(#.(A'5(P<#FC%8; %(>/5%(/<8(#<('2%(A<&"%/5A<; (#.(A'5("#C%(A<(8A5'"A>7'%8(/''/&P(8%'%&'A#<4 02%(8/5%(5'789(F/5(A@3C%@%<'%8(/5(/<(/33CA&/'A#<(#.('2%(5#=&/CC%8(D7C'A=/;%<'(695'%@)))))))))M%I%#3@%<'(bA'(ZD: 6Mb[E(8%I%#3%8(>9(/7'2#"5(#.('2A5(3/3%"4(D: 6Mb(A5(/(5#.'F/"%('##((.#"5733#"'(#.('2%('%&2<#C#; 9(#.(@7C'A=/; %<'(595'%@5(.#"@/C(53%&A.A&/'A#<(/<8(A@3C%@%<'/'A#<4(?' &#<5A5'5(#.(/(<7@>%"(#.(; %<%"A&(3"#&%87"%5E(; %<%"A&(8/'/>/5%(5'"7&'7"%5(/<8(%8A'#"5(F2A&2(@/P%(A' 3#55A>C%('#(53%&A.9(&#@3#<%<'5(#.('2%(@7C'A=/;%<'(595'%@(#<(/(.#"@/C(53%&A.A&/'A#<(C/<;7/;%(/<8 .7"'2%"( '#( ; %<%"/'%( @7C'A=/; %<'( 595'%@( /"&2A'%&'7"%E( A'5( &#@@7<A&/'A#<( &#@3#<%<'5E( '#( &C#<% P<#F(%8; %=>/5%8( /; %<'5( FA'2( '2%( <%%8%8( .7<&'A#</CA'9E( %'&4( WA57/C( 1ccE( d: W: -( /<8( eD` 5#.'F/"%(F%"%(75%8('#(8%1%C#3('2A5('##C(PA'4(\_#'A&%('2/'(8%1%C#3@%<'(#.('2%(@#8%C(/<8('2%(#>\$%&'= #"A%<'%8(3"#\$%&'(#.('2%(&/5%(5'789(F%"%('2%(@#5'('A@%(/<8(%..#"'(&#<57@A<;(3/"'(#.(F#"P(F2%"%/5(A'5 5#.'F/"%(A@3C%@%<'/'A#<(FA'2A<(D: 6Mb(8%@/<8%8(@7&2(C%55('A@%(/<8(%..#"'54 Y7"'2%"("%5%/"&2(FACC(&#<&%"<('#('2%(8%1%C#3@%<'(#.(/("A&2%"(?<37'(O"/..A&(D#8%E(%<"A&2@%<'(#. P<#FC%8; %( >/5%5( #.( 3/"'A&7C/"( /; %<'5( /<8( @%&2/<A5@5( .#"( 8%&A5A#<( @/PA<; 4( M%I %C#3@% \' ( #.( / 5#.'F/"%(.#"(#..=CA<%(/; %<'(C%/"<A<; ('#(A@3"#1%(5%&7"A'9(595'%@(3%".#"@/<&%(/<8('#(@/P%(3#55A>C9 A'5(/8/3'/>ACA'9('#(<%F(PA<85(#.(/''/&P5(A5(3C/<<%8(/5(F%CC4

# Chapter 2. Results of Research on Digital Image Steganography: Approach, Techniques, Implementation and Simulation

#### Abstract.

This Chapter outlines concisely the main steganography oriented results of the Project. These results are twofold. The first result is the development of a new approach to transparent embedding data into digital images. This approach is capable to provide a tradeoff between high rate of the embedded data and robustness to common and some intentional distortions, in particular, to JPEG compression. The developed technique makes use the properties of the singular value decomposition (SVD) of a digital image. According to these properties each singular value specifies the luminance of the SVD image layer, whereas the respective pair of singular vectors specifies image geometry. Therefore slight variations of singular values cannot influence visibly on the cover image quality. The idea of the proposed approach is to embed a bit of data through slight modifications of singular values of a small block of the segmented covers. The approach is robust because it supposes to embed extra data into low bands of covers in a distributed way. The size of small blocks is used as an attribute to achieve a tradeoff between the embedded data rate and robustness as required by the application. An advantage of the approach is that it is blind. Simulation has proved its robustness to JPEG compression up to 40%. The approach can be used both for hidden communication and watermarking.

The second result is the development of the format for compressed representation of digital images to be embedded into cover one. The idea of compression is based on using SVD of image color matrices. SVD makes it possible to represent an image as a partial sum of the most significant layers corresponding to the largest singular values. This idea is applied to each small block of the segmented image. The developed format is capable to provide compression up to 2 bit/pixel. Combined with the developed technique for hiding data in digital images, it makes possible robust embedding digital image into a cover one.

All theoretical results are validated using thorough simulation on the basis of the developed software.

The detailed description of the aforementioned results can be found in [IntRep#1], [IntRep#2] [IntRep#3] and [FinRep#1].

#### 2.1. Introduction: Overview of the Results Presented in Previous Reports.

Transparent hiding data into digital images called "digital image steganography" (DIS) presents an effective way for secret communication, watermarking and other applications. Although DIS is a quite new field of research, development of the Internet digital media and practical needs stimulated recent rapid progress in this field. Steganography by itself aims to conceal the very existence of the fact of communication. Combined with an encryption, steganography provides a higher level of the communication secrecy.

Currently this field is a subject of the intensive research. A number of techniques for DIS and watermarking have been developed during the last five years. In the Interim Report #1 [IntRep#1] the thorough overview of the state-of-the-art in DIS and watermarking areas were given. In this overview the steganography problems were analyzed in many respects. In particular, there were analyzed the terminology, classification of the embedding schemes, application-oriented classification of the embedding tasks and the respective common and particular requirements to image-based embedding techniques, general classification and overview of the developed approaches for hiding information in images. Let us summarize in brief the main results of this phase of the research.

According to the commonly accepted classification, the proposed DIS techniques can be classified as follows:

- Techniques that utilize a spatial domain. To insert data into an image, they use a selected subset of the image pixels<sup>1</sup> using a bit-wise approach ([Bender *et al-9*], [Bruyndonckx *et al-95*], [Chen *et al-99*], [Machado], [Matsui *et al-98*], [Pitas *et al-96*], [Tanaka *et al-90*], [van Schydel *et al-94*], etc.).
- Transform-based techniques, that operate with images represented by a finite set of orthogonal or bi-orthogonal functions called "basis functions" ([Burget et al], [Kundur *et al-97*], [Piva et al-97], [Podilchuck *et al-97*], [Smith *et al-96*], [Xia et al-97], [Zhu et al-95], etc.). Examples are Discrete Cosine Transform, and Wavelet transform.
- Fractal-based techniques that construct "fractal code" of an image in such a way that allows to encode both the cover and the hidden images ([Puate et al-]).

More information on this subject can be found in [Cox et al-97], [Johnson *et al-*98], [Johnson *et al-*00], [Katzenbeisser *et al-*00], [Petitcolas *et al-*99], and [Swanson *et al-*98].

The common opinion is that there is no particular superior technique. Each technique has each own merit and flaws and preferable application area.

[IntRep#1] describes detailed classification of general and particular requirements to DIS techniques. In general, the major requirements to any DIS technique are assuring invisibility of the hidden data, robustness to common and some types of intentional distortions and support a required rate of the hidden data. Since these requirements are conflicting, and concrete requirements are different for different application areas, the rational tradeoff depends on any particular application ([Petitcolas *et al-*98]).

Unfortunately, the majority of the techniques used for hiding data into digital images, known in the literature and implemented within commercial and research software tools are vulnerable to common signal processing and to intentional attacks involving distortions of sub-perceptual level. Many of the embedding systems provide a limited robustness against attacks. It was proved by employing software tools simulating such attacks, for example such tools as *StirMark*, *Unzign*, and *The Mosaic attack emulator*.

The lessons learnt by researchers on the basis of analysis of properties, advantages and disadvantages of the existing digital image steganography techniques are as follows [Petitcolas *et al-*98]:

- 1. Information hiding algorithms that attempt to meet all the accepted requirements to a steganography task solution would fail. There is no a superior solution applicable to all DIS tasks. Each solution must be based on a tradeoff depending on application, for example, robustness versus bandwidth and accessible rate of information to hide. It is quite definitely that the most real-life applications do not need to meet all requirements, to be *robust to all kinds of distortions* and to be *resistant against all* known and future types of *attacks*. For, possibly, every application one can find a technique that more or less meet the basic requirements, that more or less feasible.
- 2. Real problems are not only to insert and detect hidden data. According to [Petitcolas *et al*-98], the progress will come not just from devising new marking schemes, but in developing ways to recognize hidden data that have been embedded using obvious combinations of statistical and transform techniques and thereafter subjected to distortion.
- 3. It is common opinion that steganography would go through the same process of evolutionary development as cryptography, with an iterative process in which inventing new types of attacks will lead to the development of more robust systems [Petitcolas *et al-*98].

The above conclusions determined *the basic objectives of the steganography-oriented* research of the Project. Indeed, the study of the state-of-the-art in the DIS area exhibited that the majority of the developed techniques aims at solving the watermarking problem in which the most significant requirement is robustness to a wide range of distortions whereas high allowable rate of covertly embedded data is not an issue. That is why many existing approaches do not pay a noteworthy attention to the development of techniques capable to provide high rate of the invisibly embedded data and robustness, for example, to lossy compression like JPEG. Indeed, the only well-

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<sup>&</sup>lt;sup>1</sup> For example, using masking effect, or using a pseudo-random seeding or something other strategy.

known approach that could be used to embed an image into the cover one is embedding data into Least Significant Bit called as LSBs approach. One of such techniques was proposed in [Fridrich *et al-*99]. It uses segmentation of the image to be embedded into blocks of size  $8\times8$ , applying DCT transform to each block, quantizing and encoding its coefficients and embedding each coded block into one or into two LSBs of the cover image. Unfortunately, such and similar techniques cannot provide robustness and are highly sensitive to many common distortions, image format transformations and JPEG compression.

However, many military and industrial applications, such as hidden communication (HC) and hidden transmission of digital images call for transparent and robust embedding of high volumes of data. Examples are transmission of top-secret projects, industry secret, plans of covert operations [Johnson *et al-98*], etc. An important aspect of HC is the necessity to support the survivability of the transmitted information.

Due to the theoretical limitations it is not possible to provide both high robustness and high rate of the transparently embedded data [Anderson *et al-98*]. Nevertheless, it is highly necessary to develop an approach that should be able to provide for the high rate of the transparently embedded data preserving a reasonable robustness.

Two ideas were chosen in the development of such an approach. The first of them is the development of a new robust method for embedding data into digital images and the second one is the development a format for compressed representation of the digital image to be embedded into cover one. Both these ideas were the subjects of the research during the second and sequential phases of the work.

This research resulted in the development of a new approach to the transparent embedding data into digital images, which was described in [FinRep#1], [Gorodetski3 et al-00] and [Gorodetski4 et al-00]). The developed approach can be classified as a "Transform-based" because it deals with the image transformed to the Singular Value Decomposition (SVD). The developed method uses the properties of SVD of a digital image. According to these properties, each singular value (SV) specifies the luminance (energy) of the SVD image layer, whereas the respective pair of singular vectors specifies the image geometry. Therefore, slight variations of SVs cannot affect the visual perception of the quality of the cover image. The proposed approach is based on the embedding of data through slight modifications of SVs of a small block of the segmented cover image. The approach is potentially robust because it embeds extra data into low bands of the cover image in a distributed way. The size of small blocks can be used as an attribute to achieve a tradeoff between the embedded data rate and robustness. An additional advantage of the approach is that it is blind, i.e. it allows extracting hidden information without using of the original cover image. Below in following sections the developed approach to embedding data into digital images is described in more details.

The second idea resulted in the development of the new format for the compressed representation of digital images. This format is implemented and explored in details via simulation ([IntRep#2], [Gorodetski3 et al-00], [Gorodetski4 et al-00]). The approach is based on using singular value decomposition (SVD) of every image color matrices. SVD makes it possible to represent an image as a partial sum of the most significant layers corresponding to the largest singular values. The idea of such compression is that contribution of each i-th layer into forming of the original image is proportional to  $\lambda_i$ , since singular vectors are normalized. That is why potentially it is possible to delete the layers corresponding to the small singular values without noticeable degradation of the resulting image as compared with the original one. Extended simulation confirmed that this guesswork is valid for preserving no more than 25% of the most significant layers

The aforementioned idea was used for the development of a new image compressed format for compressed image coding. It makes use of

(1) a number of the MSLs determined according to a simple formal criterion and validated through simulation,

- (2) segmentation of the image into small blocks of a size depending on the required data bit rate to be embedded and required quality of the target image after decompression,
- (3) special quantization and
- (4) optimal encoding of singular vectors of the preserved MSLs.

Simulation proved that this format can provide less than 2 bpp data rate while preserving needed quality of the restored image. Although the last result can be thought as in some sense aside as compared with the Project objectives, nevertheless, while combining with the developed SVD-based approach for hiding data in digital images, it allows to embed robustly a digital image into a cover one. To our knowledge, there are no other methods that are capable to embed a digital image into a cover one in a way that provide robustness to the JPEG compression.

The rest of the chapter is devoted to the description of the developed SVD-based method of hiding data into digital images. In Section 2.2 the concept of the proposed approach is explained. In Section 2.3 the developed techniques of data hiding are described and the results of a simulation-based study, focused on the robustness issue as well as on embedded data rate are outlined. The developed technique is illustrated by several examples. In conclusion a general assessment of the results is given.

# 2.2. Mathematical Basis of the Developed Techniques for Hiding Data into Digital Images: Singular Value Decomposition of Digital Images

A digital image in bitmap format is specified by a  $m \times n$  matrix  $A = \{a_{i,j}\}_{m,n}$ . If an image is represented in RGB format then it is specified by three such matrices  $A_R$ ,  $A_G$  and  $A_B$ .

An arbitrary matrix A of size  $m \times n$  can be represented by its SVD ([Horn et al]) in the form

$$A = X\Lambda Y^{T} = \sum_{i=1}^{i=r} \lambda_{i} X_{i} Y_{i}^{T}$$

$$\tag{1}$$

where X, Y are orthogonal  $m \times m$  and  $n \times n$  matrices respectively,  $X_1, X_2, ..., X_m$  and  $Y_1, Y_2, ..., Y_n$  are their columns, ! is diagonal matrix with non-negative elements, and  $r \le \min\{m, n\}$  is the rank of the matrix A. Diagonal terms  $\lambda_1, \lambda_2, ..., \lambda_r$  of matrix ! are called  $singular \ values$  (SV) of the matrix A and r is the total number of non-zero singular values. Columns of the matrices X, Y are called left and right singular vectors of the matrix A respectively. Singular values  $\lambda_1, \lambda_2, ..., \lambda_r$  can be calculated as  $\lambda_i = \sqrt{\mu_i}$  i=1,2,...,r, where  $\mu_i$  is an eigenvalue of the matrix  $AA^T$ , or  $A^TA$ . The left singular vector  $X_i$ , i=1,2,...,r, is equal to the eigenvector of the matrix  $AA^T$  corresponding to  $\mu_i$ . Similarly, the right singular vector  $Y_i$ , i=1,2,...,r, is equal to the eigenvector of the matrix  $A^TA$  that corresponds to its eigenvalue  $\mu_i$ . If an image is given in RGB format then it is represented by three SVDs in the form (1).

Thus, SVD of an image decomposes the respective matrix into layers  $\lambda_1 X_1 Y_1^T$ ,  $\lambda_2 X_2 Y_2^T$ ,...,  $\lambda_s X_r Y_r^T$ . As a rule, SVs are enumerated in descending mode, i.e. if  $\lambda_i > \lambda_j$  then i < j, and  $\lambda_1$  is the maximal SV.

SVD possesses several interesting properties ([Horn *et al*]; two of them are utilized below to achieve invisible and robust hiding of extra data in digital images.

The first property is that each SV specifies the luminance (energy) of the SVD image layer, whereas the respective pair of singular vectors specifies an image "geometry". It was discovered that slight variations of SVs do not affect visual perception of the quality of the cover image. This property is used to embed a bit of data through slight modifications of SVs of a small block of a segmented cover.

The second property is that without the loss of image quality an image could be represented by so-called Truncated SVD (TSVD), i.e. by the sum

$$A_s = \sum_{i=1}^{i=s} \lambda_i X_i Y_i^T, s << r.$$
 (2)

instead of sum (1) ([Gorodetski3 et al-00]). In other words, a TSVD of an image can be used for its compressed representation. SVD-based image compression was proposed in [Andrews et al-76]. Later a number of approaches that combine SVD image transform with other transforms was developed. However, the major attention was paid to the development of a lossyless SVD image compression and its combinations with other transforms. For example, to code images, in [Yang et al-95] SVD transform is combined with Vector Quantization approach, in [Waldemar et al-97] a combination of SVD and Karhunen-Loeve transform is used to develop a hybrid compression. In [Fukutomi et al-99] SVD transform is combined with wavelet transform. In [Gorodetski3 et al-00] a format for SVD-based lossy compression of digital images was proposed. This format provides the rate of compression close to 2 bit/pixel while preserving the appropriate quality of the restored image. This format is used to solve the task of robust embedding a digital image into a cover one.

#### 2.3. SVD-based Techniques for Hiding Data into Digital Images

The following are the techniques that were developed to utilize the first of the two aforementioned properties of SVD image representation.

#### **2.3.1. Technique 1**

In brief, the first proposed technique is as follows. A cover image represented in 24 bpp (RGB) format is segmented into blocks of size s×s¹ and SVDs for each such a block and for each matrix of Red, Green and Blue layers are computed. Each block of every color layer is used to embed a bit of data. In the Technique 1, a bit of data is embedded through a slight modification of the largest singular value of the block. The implemented and explored algorithm of modification is described below.

Let B(k,l) be a block, where k is the block number and  $l \in \{Red, Green, Blue\}$ . Let the largest SV of the block B(k,l) of size s×s be  $\lambda_1^k$ . Let b be a bit of data to be embedded into this block. The embedding algorithm is as follows:

For each pair (k, l)

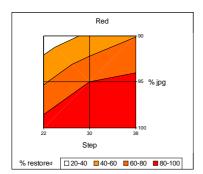
- 1. Choose the quantization step, d, of the largest singular value of the block. The value d may be different depending on the layer of color.<sup>2</sup>
- 2. Compute integer number S such that  $\lambda_1^k(l) = S \times d + \delta$ ,  $\delta < S$ .
- 3. Embed bit *b* of data as follows:

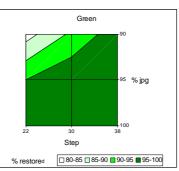
If S is the odd number then
if b=1 then S is not changed
else
if b=0 then S:=S+1.

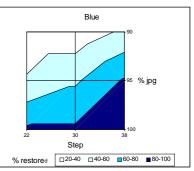
If S is the even number then
if b=1, then S:=S+1, else
if b=0, then S is not changed.

<sup>1</sup> The developed software implements a particular case of this technique when s=4, generally, this number could be extended. Note that an increase of k value results in the increased robustness of the technique and in the decreased volume of transparently embedded data.

<sup>&</sup>lt;sup>2</sup> The value of *d* is selected on the basis of statistical exploration of correlation between distortion caused by JREG compression of various percentages and the probability of a bit recovery. This exploration must be made for each color. The respective results are given below in this section.







**Fig. 2.1.** The plots of dependencies between the step of quantization *d* of the largest singular value of a cover image and survivability of the embedded image after JPEG compression for *red* (left), *green* (center) and *blue* (right) color layers of the cover image. These results correspond to the case when cover image is segmented into blocks of size 4×4

4. Compute the modified singular value  $\widetilde{\lambda}_1^k(l)$ :

$$\widetilde{\lambda}_{1}^{k}(l)=d\times S+d/2.$$

5. Compute the matrix of the block having modified largest singular value:

$$\widetilde{B}(k,l) = \widetilde{\lambda}_{1}^{k}(l)X_{1}(l)Y_{1}^{T}(l) + \sum_{i=2}^{4} \lambda_{i}(l)X_{i}(l)Y_{i}^{T}(l)$$

6. Result: Matrix  $\widetilde{B}(k,l)$  containing embedded bit of data.

The major implementation concern of the above algorithm is how to choose the quantization step d. It is obvious that the increase of the value d leads both to more robust data hiding and to less transparency of the embedded data. To explore this dependency quantitatively, the respective simulation was performed for several cover images segmented into blocks of size 4×4 pixels and several images to be embedded. The latter were segmented into small blocks of size 8×8 and then compressed using TSVD proposed in [Gorodetski3 et al-00] and mentioned in Section 2. The results are given in Fig.2.1a (for Red layer), Fig.2.1.b (for Green layer) and Fig.2.1.c (for Blue layer). The plots illustrate the correlation between the percentage of JPEG compression and percentage of the correctly recovered blocks of the embedded image. Based on this result the following values of quantization steps within the Red, Green and Blue layers were chosen:

$$d(red)=46$$
,  $d(green)=22$  and  $d(blue)=52$ .

It should be noticed that the quantization step value *d* might be used as a component of the secret key providing the restricted access to the hidden information.

The simulation indicated that this way of embedding data into a cover image is robust against 100% JPEG compression.

The hidden data extraction procedure is very simple. Let  $\widetilde{B}(k,l)$  be a block with an embedded bit of data.

For each pair (k, l) do

- 1. Compute the largest singular value  $\widetilde{\lambda}_1^k(l)$ .
- 2. Compute  $\widetilde{\lambda}_1^k(l)/d=S+d/2$ .
- 3. If *S* is even number then the embedded bit value is *0* otherwise it is *1*.

<sup>&</sup>lt;sup>1</sup> Note that the appropriate choice is specific for every way of cover images segmentation.



**Fig.2.2.a.** Cover image. It is presented in RGB format and is of size 600×512 pixels



**Fig.2.2.b.** Image to be transmitted. It is gray and is of size 240×120



**Fig.2.2.c.** Recovered hidden image. The stego-image was subjected to JPEG compression

Fig. 2.2. An example of the use of the Technique 1 for embedding image into cover one

This approach to data embedding provides a sufficiently high rate of the embedded data although its bit rate is less than the one provided by LSBs techniques. For example, let cover image be of size  $600\times512$  (see Fig.2.2.a) and s=4. It comprises  $150\times128$  blocks of size  $4\times4$  in each color. Therefore this technique makes it possible to embed up to 57600 bits. The picture to be embedded (see Fig.2.2.b) is of size  $240\times120$  and segmented into 200 blocks of size  $12\times12$ . Due to TSVD compression (see Section 2.2), each such block is represented into blocks of  $12\times12$  using segmentation. Its length is 288 bits, that is why the total size of the image of Fig.2b in TSVD format is equal to 56000 bits. Hence, it is possible to embed it into image depicted in Fig.2.2.a.

This technique was subjected to several experimental studies. An example of such a study is given in Fig.2.2. The cover image (Fig.2.2a) containing embedded image (Fig.2.2b) was distorted by JPEG compression and then transformed back in BMP format. The hidden image extracted from the JPEG distorted stego-image is depicted in Fig.2.2.c. One can see that the quality of the reconstructed image is not excellent but is very satisfactory. Notice that the bit rate of the TSVD image to be embedded can be increased using a Haffman-like compression of the TSVD files for each block.

Embedding and recovery procedures could be equipped with a secret key to seed blocks thus providing additional level of security. To improve the survivability of the hidden data during transmission it is possible to transmit the same image several times using various covers. It should be noticed that if the image to be hidden is of very large size then its transmission could be implemented using several cover images.

#### **2.3.2. Technique 2**

This technique uses a different approach to embed data into a cover image. Notice that the type of embedded binary file is irrelevant. A bit of information is embedded into a block of the segmented cover image. The block size could be chosen arbitrarily.

Let cover image be represented in a 24 bit (RGB) format. Data can be embedded independently into each RGB layer of the cover image or, optionally, in specified layer(s). Let the size of blocks of the cover image segmentation be  $m \times k$  pixels, A be the matrix of a block of the covers corresponding to a color layer from Red, Green or Blue. Let a bit b to be embedded into block A of a layer.

The algorithm of the *embedding procedure* is as follows:

- 1. Compute singular value decomposition of matrix A. Let  $V^{\lambda} = [\lambda_1, \lambda_2, ..., \lambda_r]$  be the vector of singular values of matrix A ordered in decreasing mode,  $X_i$  and  $Y_i$  are singular vectors of matrix A and i=1,2,...,r, where r is the rank of matrix A.
- 2. Compute Euclidean norm of vector  $V^{\lambda}$ ,  $Norm(V^{\lambda}) = \sqrt{\sum_{i=1}^{r} (v_{i}^{\lambda})^{2}}$ , where  $v_{i}^{\lambda}$ , i=1,2,...,r, are the components of vector  $V^{\lambda}$ .
- 3. Select the value of Delta that is the step of quantization of the Euclidean norm of vector  $V^{\lambda}$ . Remark 1. The appropriate value of Delta depends on the color layer used for embedding a bit of data and has been chosen for each layer through simulation. Notice that the value of Delta can play the role of a component of the secret key restricting access to the hidden data. One more way to increase the secrey is to use an uneven quantization of  $Norm(V^{\lambda})$ .
- 4. Compute the integer  $N=[Norm(V^{\lambda})/Delta]$ , where [\*] is the integer part of the quotient of division.
- 5. Embed bit *b* according to the following algorithm:

If 
$$b=1$$
 then   
{if N is odd then  $\tilde{N}=N+1$  else  $\tilde{N}=N$ } else (if  $b=0$ )   
{if N is even then  $\tilde{N}=N$  else  $\tilde{N}=N+1$ }}

6. Compute the modified norm of the vector of the singular values:

$$Norm(\widetilde{V}^{\lambda}) = \widetilde{N} \times Delta + (Delta/2).$$

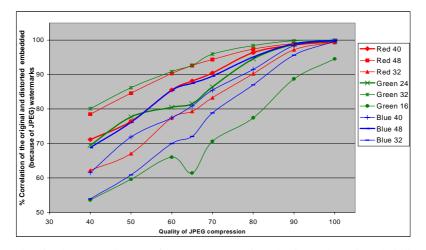
7. Compute the modified vector of the singular values:

$$\widetilde{V}^{\lambda} = V^{\lambda} \times (Norm(\widetilde{V}^{\lambda})/Norm(V^{\lambda})).$$

8. Compute the modified matrix of the block in which bit b is embedded:

$$\widetilde{A} = \sum_{i=1}^{r} \widetilde{\lambda}_{i} X_{i} Y_{i}^{T}.$$

9. End of the embedding procedure.



**Fig. 2.3.** Dependencies between degree of JPEG compression (horizontal) and probability of the watermark presence (vertical) for various value *Delta*. (Averaged over a set of images)

The described algorithms must be applied to each block of the covers in which a bit of data is to be embedded.

The embedding procedure may provide restricted access to the hidden data via using seeding of the binary string to be embedded in pseudo random mode or via transpositions of the lines and columns [Gorodetski3 *et al-*00].

The extraction task is simpler then the embedding. Let  $\widetilde{A}$  be the matrix of block containing hidden bit b of data, which must be extracted.

- 1. Compute the singular value decomposition of block  $\widetilde{A}$ . Let  $\widetilde{V}^{\lambda}$  be vector of singular values ordered in decreasing fashion,  $X_i$  and  $Y_i$  are singular vectors of matrix  $\widetilde{A}$  and i=1,2,...,r, where r is the rank of matrix  $\widetilde{A}$ .
  - Remark 2. SVs and singular vectors of the same block  $\vec{A}$  of the modified cover image can be additionally modified during transmission or changed intentionally. For simplicity, we denote them in the same way as they were denoted in the embedding procedure.
- 2. Compute the Euclidean norm of vector  $\widetilde{V}^{\lambda}$ , i.e. Norm  $\widetilde{V}^{\lambda} = \sqrt{\sum_{i=1}^{r} (\widetilde{V}_{i}^{\lambda})^{2}}$ .
- 3. Compute  $\tilde{N} = [Norm\tilde{V}^{\lambda}/Delta]$ , where [\*] is the integer part of the quotient of division.
- 4. Compute the value of the hidden bit *b*:

{If 
$$\widetilde{N}$$
 is even then  $b=1$  else  $b=0$ }

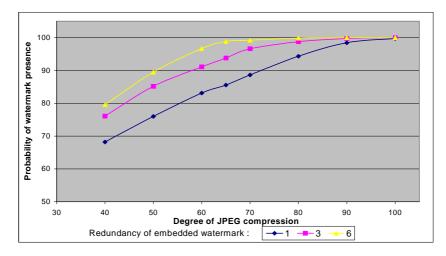
5. End of extraction procedure.

The described embedding (extraction) algorithm must be applied to each block of the covers (stego-image), in which a bit of data is embedded. The extracting procedure may require the knowledge of the secret key if it has been applied in embedding.

This technique was investigated statistically from several points of view. The first task was to explore the optimality of *Delta* values in the trade-off between the robustness of the hidden data and its visibility. It was established that the appropriate value of *Delta* depends on the color layer. In fig.3 the results of the simulation-based investigation of the aforementioned dependencies are given. One can see that optimal choices of the of Delta given color layer are close to the followings:

Given such values of *Delta*, Technique 2 proves robustness to JPEG compression up to degree 40% provided that the value of the watermark presence probability is (0.7–0.8).

The special attention was paid to study the relationship between the degree of JPEG,



**Fig. 2.4.** Dependency between *Degree of JPEG compression* (horizontal axis) and *Probability of the watermark presence* for various degrees of redundancy

compression redundancy of the embedded watermark and the probability of the watermark presence. In this study redundancy is understood as the number of embedded copies of the watermark. The results are displayed in Fig.2.4. One can see that the redundancy presents an additional way to increase the robustness of Technique 2.

In fig.2.5 an example of using Technique 2 for embedding emf-file into a digital image is given. Left-hand image (fig.5a) corresponds to the cover image (in gray scale) with the embedded emf-file. The right-hand image depicts the cover image with extracted data that indicates, for example, to a pilot the route to follow that was hidden into transmitted image.

#### 2.4. Conclusion of the Chapter 2

This Chapter presents the main final result of the steganography-oriented research on the Project. It is a novel approach to digital image steganography and two particular techniques implementing this approach. Both techniques utilize the concept of embedding data through slight modifications of singular values of small blocks of the cover image.

The techniques, implemented in software, are subjected to statistical analysis of their robustness and the allowable rate of covertly embedded data. In particular, it is shown the developed approach for data hiding is robust to JPEG compression up to 40%. It is blind, i.e. the hidden data can be extracted without possessing the original cover image. Embedding and recovery procedures can be equipped with a password and secret key to seed blocks thus providing restricted assesses to the hidden data. There exist several additional attributes resulting in additional security assurance.





**Fig.2.5.a.** Cover image containing invisibly embedded *emf*-file representing an aircraft route

**Fig.2.5.b.** Extracted route drawn over the cover image

**Fig. 2.5.** An example of the use of the Technique 2 to hide route of an aircraft into the map. The hidden data are represented in *emf* graphical format.

The analyses indicate that the developed techniques are suitable for hidden communication that calls for high rate and appropriate survivability of the embedded data, and invisibility of the hidden image, and for watermarking where the main requirement is robustness to common and some intentional distortions.

The approach is demonstrated by several examples of practical applications.

### 3. General Conclusion on the Project

According to the Project Work Plan the main tasks that the Project was addressing to are as follows:

- 1) Development of the architecture of the agent-based information security system on the whole and architectures of particular agents; development of the ontology of information security domain to design and to decompose distributed knowledge base structure.
- 2) Development of a formal framework for representation of the agents' distributed knowledge.
- 3) Development of the procedure of the agents' cooperation for integrated information security task solving.
- 4) Development and mathematical justification of the new methods of image-based information hiding (image-based steganography) to provide safe channels of information exchange.

All aforementioned tasks were solved and the results were submitted to the Partner in three Interim and two Final reports. The results of the research were published in proceedings of seven international conferences and workshops with reviewed articles. Two papers have been accepted already to the forthcoming international workshop on computer network security.

This Report presents the final results of the research. The Project objectives include two in some sense independent tasks. One of them concerns to the multi-agent model and software prototype of the computer network security system and the other one concerns to the digital image steganography. Accordingly, this Final Report comprises two chapters. Each of them is devoted to the brief outline of the respective results submitted in previous reports and to the description of the new results that were achieved during the final phase of the research. In both Chapters we accentuate the software developed for verification and validation of the theoretical results, i.e. for verification and validation of the developed approaches, methods, architectures, techniques and algorithms associated with the Project tasks.

The Chapter 1 considers the developed Case Study of the network security system that is a software implementation of the multi-agent security system. The Case Study is composed of particular autonomous knowledge-based agents, distributed over the hosts of the computer network to be protected and cooperating to make integrated consistent decisions. The Chapter 1 describes the architecture of the security system Case Study and architectures of its components that are particular software agents, communication components and software that is intended to simulate the input traffic. The Case Study architecture corresponds to the multi-agent system protecting a segment of a Local Area Network comprising four hosts to be protected. Each host-based component of the network security system comprises seven specialized software agents situated on the host. In total, the Case Study comp[rises 28 distributed over the network interacting software agents.

Simulation of the Case Study displayed a number of advantages of multi-agent architecture for protection of a computer network against distributed attacks. The most significant of them is a capability of the comparatively "light" components of a multi-agent security system to cooperate to solve a "heavy" task. At present the only way to detect efficiently a distributed attack against a computer network is a cooperation of security software entities (agents) distributed over the hosts of the network and within each host itself. A sample of usefulness of agent cooperation can be seen within the described example. In particular, detection of a *combined spoofing attack* has become possible only due to cooperation of the software agents situated on different hosts. One more example of the necessity and usefulness of the cooperation of security agents is operation of the knowledge-based agent named IDA2. It is intended to collect information about suspicious behavior of users at many entry points of the native host and exchange information with similar agents of the other hosts to make integral decision about status of connections. In the developed case study this agent has been provided by a comparatively poor knowledge base and can not play a significant role in intrusion detection. In the further development of the case study a significant

accent must be made on enrichment of its knowledge base and on the increasing its role in distributed attack detection.

This Case Study is implemented as distributed multi-agent system, which components interact via message exchange. Simulation scenario, input traffic model and peculiarities of the distributed security system operation are described. The major attention is paid to the intrusion detection task and agents' interactions during detection of an attack against the computer network. Case Study implementation was carried out on the basis of the Multi-agent System Development Kit developed by authors of the research.. The software code is developed using Visual C++, JAVA 2 and XML.

The Chapter 2 outlines concisely the main steganography oriented results of the Project. These results are twofold. The first result is the development of a new approach to transparent embedding data into digital images. This approach is capable to provide a tradeoff between high rate of the embedded data and robustness to common and some intentional distortions, in particular, to JPEG compression. The developed technique makes use the properties of the singular value decomposition (SVD) of a digital image. According to these properties each singular value specifies the luminance of the SVD image layer, whereas the respective pair of singular vectors specifies image geometry. Therefore slight variations of singular values cannot influence visibly on the cover image quality. The idea of the proposed approach is to embed a bit of data through slight modifications of singular values of a small block of the segmented covers. The approach is robust because it supposes to embed extra data into low bands of covers in a distributed way. The size of small blocks is used as an attribute to achieve a tradeoff between the embedded data rate and robustness as required by the application. An advantage of the approach is that it is blind. Simulation has proved its robustness to JPEG compression up to 40%. The approach can be used both for hidden communication and watermarking.

The second result is the development of the format for compressed representation of digital images to be embedded into cover one. The idea of compression is based on using SVD of image color matrices. SVD makes it possible to represent an image as a partial sum of the most significant layers corresponding to the largest singular values. This idea is applied to each small block of the segmented image. The developed format is capable to provide compression up to 2 bit/pixel. Combined with the developed technique for hiding data in digital images, it makes possible robust embedding digital image into a cover one.

All theoretical results of the Project are validated using thorough simulation on the basis of the developed software. The developed software will be demonstrated and submitted to the Partner.

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### Appendix A1. Case Study Model

### A1.1. Structure of the basic concepts of the case study

One of the main concepts of the case study is the concept of "current connection". This concept is for formation in CNSS of the generalized representation of the input messages (packets) sequence of one connection, and also for reflection of a current status of connection on the basis of processing of the messages, which has arrived to a current moment. This concept is a base for determination of the CNSS performance scheme. The CNSS performance process can be represented as the distributed processing of examples of this concept by various classes of the security agents according to their specialization.

When the first message, intended for a new connection activating, arrives on a host the agent AD-E of this host forms the description of the connection as an example of this concept. The attribute values of this example can vary dynamically as a result of the subsequent messages processing within the current connection. As the process of connections processing in a multiagent system is distributed, the different agents can have information on the same connection simultaneously. This information processing is coordinated according to the set scheme of agents' interaction. At connection completion an appropriate example of concept "the current connection" is transferred in the completed connections archive.

The list of attributes used in the model of concept "current connection", which is operated with the security agents of a host, is presented in the table A.1.

Table A.1. The list of attributes used in model of concept "current connection"

ID	Unique identifier of an example of the concept "current connection"	
IP address	IP-address of a host establishing the connection	
Port destination	Number of a port, with which the connection is installed	
Flag	Connection status	
Last sqn_no	Sequence number of the last obtained tcp-packet	
Access	Access rights	
Time on	Time of the connection beginning	
Time last tcp	Time of arrival of the last message obtained in connection	
Time life	Life time of the connection in a "current status"	
Fact	Fact of suspicious action, detected in connection	
Data	Contents of a field "Data" of the last tcp-packet obtained in connection	
N	Reserve numerical field	

The value of a field *ID* is an unambiguous identifier of the example of concept "current connection".

The values of attributes *IP address, Port destination* and *Time on* are assigned on the basis of parameters of the first *tcp*-message in connection and do not vary.

The field *Flag* identifies a current status of the connection. The list of possible statuses of the connection is defined on the scheme of transitions of the current connection statuses (Fig.A.1). *The main states of the connection* are the following:

- HOC (Half Open Connection) status originating after the first phase of "hand shake",
- Set connection status originating after the third phase of "hand shake",
- Login status originating after identification and authentication of the user,
- *Closed* completed connection (a transition in this status is carried out after generation of the TCP-message with a flag Fin or on time).

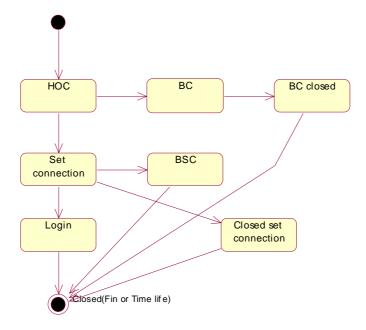


Fig.A.1. Scheme of transitions of the current connection statuses

The variation of the enumerated statuses of the connection happens as a result of arrival of the new messages.

Besides the statuses of the connection can vary automatically after defined time intervals. Thus the following classes of the connection statuses are selected:

- BC (Bad connection) the status, in which the connection transfers from a status HOC after a time interval Ack\_wait\_time,
- BC closed the status, in which the connection transfers from a status BC after a time interval Bad\_connection\_life,
- BSC (Bad set connection) the status, in which the connection transfers from a status Set connection after a time interval Set connection life.

The time intervals of the connection existence under absence of the new messages are determined by means of values of attribute *Time life*. The instant defined by attribute *Time of last tcp* is used for time intervals counting. Besides the enumerated classes of the connection statuses defined as a result of the messages obtaining or on time, one more intermediate status causing suspicions is identified:

• *Closed set connection* - status originating after obtaining of the message with a flag Fin, when the connection is in a status *Set connection*.

The value of attribute *Last sqn\_no* is updated at arrival of the new message. This attribute value is a formal tag for determination of the following message in connection.

The value of attribute *Access* sets within the connection access rights, which are determined during identification and authentication of the user.

If during the connection an action suspicious from a security view point is found out, this fact is fixed with usage of the attribute *Fact* value. If in the connection some facts of suspicious actions were detected, then a name of the last detected action class corresponds to this attribute. Thus there is no "loss" of the facts detected earlier, as the connections with the detected separate suspicious facts are recorded in the database of the agent *IDA2*. Thus, if during the connection some facts of suspicious actions are detected, all of them are recorded in the agent *IDA2* database.

In a field *Data* the data string from the same field of the last obtained message is stored.

Thus, the concept "current connection" is a basic concept for organization of the agents' interaction in CNSS and for generation of the contents of the transmitted messages. On the basis of this concept the number of additional concepts for detection of the attack classes "Denial of

service" and "port scanning" is implemented. The agents *AD-P1* use these concepts, therefore the determination of the concepts is considered at the description of the scripts of this class agent behavior.

# A1.2. Scripts of the agents' behavior

In this subsection the behavior scripts of the following agents' classes is determined:

- AD-E agents demons of preprocessing of the input tcp-messages,
- *AD-P1* agents demons of pattern detection in connection till the moment of the users' authentication and identification,
- *AD-P2* agents demons of pattern detection in connection after the moment of the users' authentication and identification,
- AIA identification and authentication agents,
- ACA access control agents,
- IDA1 intrusion detection agents revealing a combined spoofing attack class,
- *IDA2* intrusion detection agents revealing the complex attacks scripts.

The description of the all classes agents' behavior is fulfilled under the following scheme.

At first a diagram is represented which determines (1) agents' classes with which the agent of the circumscribed class cooperate, (2) content of the input and output messages as concept classes, (3) enumeration of the agents' behavior scripts defined by the meta-rules. If for the agents'

class the more than one behavior script is determined, then a common meta-script of the agent's behavior is also defined. This meta-script is responsible for choice of the necessary agent's behavior script depending on the class and content of the input message.

Then a more detailed description of each behavior script is defined. The input of the scripts is fulfilled by means of the special editor that is a MASDK component. The user interface of this editor consists of several windows. The main window of the editor maps a general behavior script as a decision tree. The auxiliary windows map a detailed representation of separate rules of a decision tree. Further in the appendix text a common definition of the agents' behavior scripts is given.

The Fig.A.2 (where MR – meta-rule) allows to restore a correspondence between the earlier circumscribed general representation of the behavior scripts (subsection 1.4) and that representation which is used in the main window of the editor.

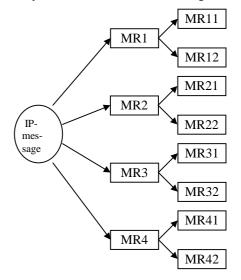


Fig.A.2. The representation of the agent *AD-E* behavior script (Script "A")

### A1.2.1. Scripts of the agent AD-E operation

The main task of an *agent-demon* is the preprocessing of the input tcp-messages arriving from *Input Traffic Model* (Fig.A.3). The preprocessing of such messages is fulfilled using data about the current connection, to which the input message concerns. Besides this messages the marks of model time and the facts concerning to current connections in a format of the current connection description can arrive on an input of the agent. The agent's behavior script is selected depending on the contents of the arrived message on the basis of preset behavior model: the script A is selected for processing the tcp-messages, the script B - for processing the time stamps *time*, the script C - for processing the facts concerning to current connections.

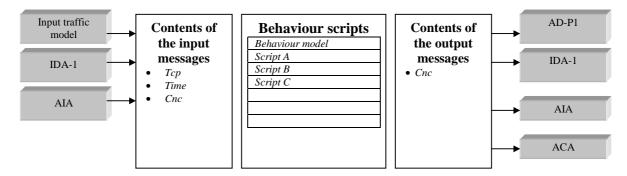


Fig.A.3. Common scheme of the agent *AD-E* operation

## Script A. Common scheme of preprocessing of the obtained tcp-messages

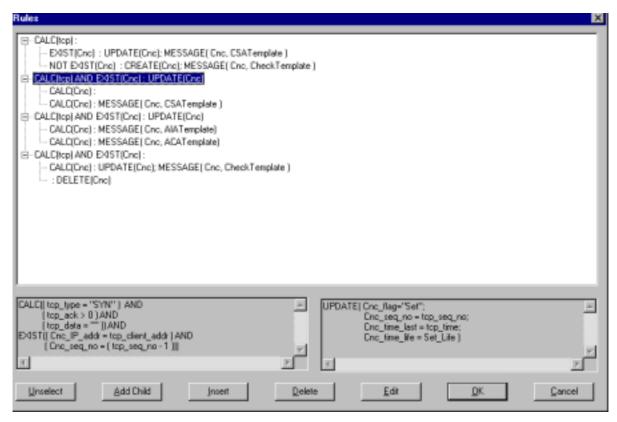


Fig.A4. Rules of the input tcp-messages preprocessing

Fig.A.4 represents four generalized rules of the agents' class *AD-E* behavior (see upper window).

A.1. The rule determines that the input tcp-message with a flag Syn is obtained, this message corresponds to a new connection establishment (the generalized condition of this rule is shown in Fig.A.4, line 1 in the upper window). According to this rule two decision variants are realized depending on various conditions.

A.1.1. The input tcp-message with a flag Syn is obtained, this message corresponds to a new connection establishment from the part of some host H, concerning to a defended network (generalized condition and action of this rule are shown in Fig.A.4, line 2 in the upper window). The *agent - demon* has information that the attack "Denial of service" is accomplished on the

host H. The record in the form "current connection" in this agent's database (see table A.2) testifies this fact. This record was registered in an agent's database during execution of the behaviour rule C.1 of the same agent (see below). As a result of the obtained message processing a modification of the available record about current connection is fulfilled (see table A.2), and this information is addressed to the agent *IDA1* on the same host.

Table A.2.

	Initial record	Modification of the record attributes
ID	-	Current N + 1
IP address	Tcp ( Client.address )	=
Port destination	-	Tcp ( Server.port )
Flag	-	НОС
Last sqn_no	-	Tcp ( sqn_no )
Access	-	-
Time on	-	Tcp ( time )
Time last tcp	-	Tcp ( time )
Time life	-	Ack_wait_time
Fact	Denial of service	Connect imitation
Data	-	-

Designations: - - the value is not defined, = - a former value.

A.1.2. The obtained input tcp-message with a flag Syn corresponds to a new connection establishment (see Fig.A.4, line 3 in the upper window). The registration record of a new connection (table A.3) is generated in the agent's database in a format of the concept "Current connection", and this information is transferred to the agent *AD-P1*.

Table A.3.

ID	Current N + 1
IP address	Tcp ( Client.address )
Port destination	Tcp ( Server.port )
Flag	HOC
Last sqn_no	Tcp ( sqn_no )
Access	-
Time on	Tcp (time)
Time last tcp	Tcp ( time )
Time life	Ack_wait_time
Fact	-
Data	-

A.2. The rule determines, that the input tcp-message with a flag Ack is obtained, this message corresponds to the last phase of "hand shake" (see Fig.A.4, line 3 in the upper window and also contents of two windows below, this rule is marked). On the basis of a sequence number a record appropriate to this connection is found in a database, and on the basis of parameters of the obtained message a modification of this record is fulfilled (table A.4).

Table A.4.

	Modification of the record attributes
ID	=
IP address	
Port destination	=
Flag	Set connection
Last sqn_no	sqn_no + 1
Access	-
Time on	=
Time last tcp	Tcp ( time )
Time life	Connection life time
Fact	=
Data	-

- A.2.1. The field *Fact* of current connection does not contain any data (Fig.A.4, line 5 in the upper window). This condition corresponds to normal connection.
- A.2.2. The value of the field *Fact* of current connection is Denial of service (Fig.A.4, line 6 in the upper window). In this case the value of the field *Fact* varies on *Connect imitation*, and this information is transferred to the agent *IDA1*.
- A.3. The rule determines that the tcp-message transmitting the data is obtained (see a Fig.A.4, line 7 in the upper window). On the basis of the message sequence number in a data base a record appropriate to this connection is found, and on the basis of the obtained message parameters a modification of this record attributes (table A.5) is fulfilled.

Table A.5.

	Modification of the record attributes
ID	
IP address	
Port destination	
Flag	
Last sqn_no	Tcp ( sqn_no )
Access	
Time on	
Time last tcp	Tcp ( time )
Time life	Connection life time
Fact	
Data	Tcp (data)

A.3.1. The rule determines that the connection is in a status *Set connection*, and dispatches to agent *AIA* the message about the connection, including data from a field data of the input *tcp*-message (see Fig.A.4, line 8 in the upper window).

- A.3.2. The rule determines that the connection is in a status *Login*. Besides the rule dispatches to agent *ACA* the message about the connection, including data from a field data of the input *tcp*-message (see Fig.A.4, line 9 in the upper window).
- A.4. The rule determines that the *tcp*-message with a flag Fin is obtained, this message corresponds to the connection closing (see Fig.A.4, line 10 in the upper window). On the basis of the message sequence number a record appropriate to the connection is found in a database.
  - A.4.1. The rule determines that the connection is in a status *Set connection* (see Fig.A.4, line 11 in the upper window). In this case the value *Empty connection* is assigned to attribute *Fact* in a record about the connection and the information about a connection is transferred to the agent *AD-P1*.
  - A.4.2. The record about the connection is deleted from the agent's database (see Fig.A.4, line 12 in the upper window). If the description of the connection contains any detected fact of non-authorized actions, then the information on the completed connection is recorded in archive of the completed connections.

Script B. Supervision of a life time of the connections in current statuses

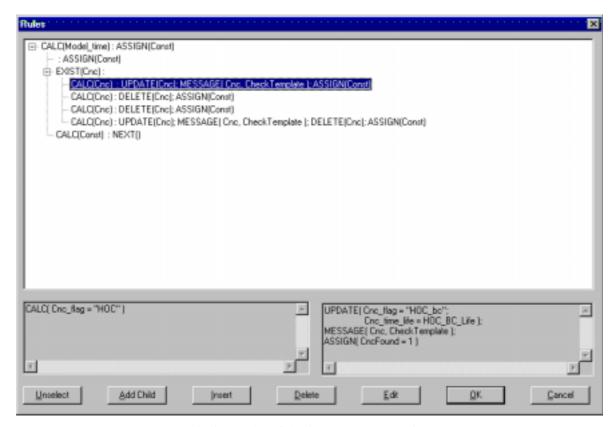


Fig.A.5. Rules of the time stamps processing

- B.1. The condition of the rule determines that the message on a current value of the model time is obtained (see Fig.A.5, line 1 in the upper window).
  - B.1.1. The rule sets the order of execution of the subsequent rules of this level in a decision tree (see Fig.A.5, line 2 in the upper window).
  - B.1.2. The rule discovers the first record about the connection, in which the connection life time is delayed in a current status (a selection condition is *Exist Connect*, *where time\_last\_tcp + life\_time > current\_time*) (see Fig.A.5, line 3 in the upper window).

- B.1.2.1. The connection from a status *HOC* (*Half open connection*) is transferred in a status *Bad connection*, the appropriate life time value is assigned to the attribute *Life\_time*, and the message on the connection status is transferred to the agent *AD-P1* (see Fig.A.5, line 4 in the upper window and also contents of two windows bellow, this rule is marked).
- B.1.2.2. The connection from a status *Set connection* is transferred to the status *Bad set connection* and this information is transmitted to the agent *AD-P1*, then the record about the connection is deleted from a database and is recorded in archive of the completed connections (see Fig.A5, line 5 in the upper window).
- B.1.2.3. The connection is in a status *Log*. In this case a record is deleted from a database (see Fig.A.5, line 6 in the upper window).
- B.1.2.4. The connection from a status *Set connection* is transferred to *BSC* (*Bad set connection*) and this information is transmitted to the agent *AD-P1*, then the record about the connection is deleted from a database and is recorded in archive of the completed connections (see Fig.A.5, line 7 in the upper window).

Script C. Responses on the facts about the connections

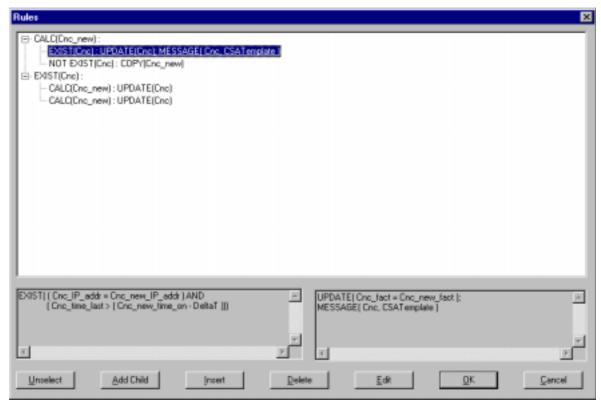


Fig.A.6. Rules of the registration of the detected facts in connections

- C.1. The rule determines that the message (from the agent *AD-P1* of other host) about detection of attack Denial of Service has come (see Fig.A.6, line 1 in the upper window). The message arrives in the format of table A.6.
  - C.1.1. The rule searches a record about the connection from a host H opened in a preset time interval (see Fig.A.6, line 2 in the upper window and also contents of two windows below, this rule is marked). The search of such record is fulfilled using the address of a host H and the defined time interval. If such connection is retrieved, then in appropriate record the value *Denial of service* of attribute *Fact* varies on a value *Connect imitation*.

C.1.2. The rule checks that during a preset time interval the connection from a host H was not open (see Fig.A.6, line 3 in the upper window). In this case the obtained warning is recorded in the agent's database and is used further during the rule A.1.1 execution.

Table A.6

	Initial record
ID	-
IP address	Tcp ( Client.address ) Host H
Port destination	-
Flag	-
Last sqn_no	-
Access	-
Time on	-
Time last tcp	-
Time life	-
Fact	Denial of service
Data	-

- C.2. The rule determines that the answer message from the agent AIA about the identification and authentication results has come (see Fig.A.6, line 4 in the upper window).
  - C.2.1. The value of a flag *Log* in the obtained connection testifies to a successful identification and authentication (see Fig.A.6, line 5 in the upper window). In this case in the connection record of the agent's database the status flag changes to *Log*, and the access level defined by the agent *AIA* according to access control rules is assigned to attribute *Access*.
  - C.2.2. This case corresponds to a negative result of the identification and authentication (see Fig.A.6, line 6 in the upper window). The connection remains in the status *Set connection*, and the value *Bad login* is assigned to the attribute *Fact*.

# A1.2.2. Scripts of the agent AD-P1 operation

The main function of this class agent (see Fig.A.7) is a support of a current quantity statistics of open connections for detection of the *port scanning* and *Syn flood* attack of the class *Denial of service*. Fig.A.8 can be an informative explanation of the behaviour script, defined for this agent. Formal tag of port scanning is the quantity of the connections opened from one host. The

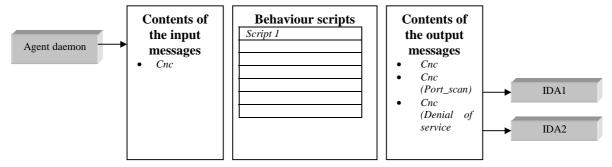
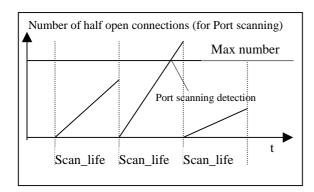


Fig.A.7. Common scheme of the agent AD-P1 operation

conclusion about possible port scanning can be done if the connections quantity will exceed a preset threshold. This conclusion is for construction of a rule of the port scanning detection. However on this basis it is possible to generate a more exact rule taking into account variety of



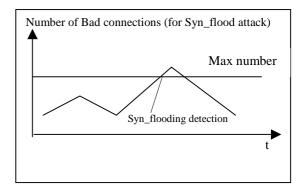


Fig.A.8. Graphical representation of rules of attack detection by the agent AD-P1

ports and also presence of informative actions in these connections.

A formal tag of the *Syn flood* attack realization is a quantity of connections on the same port in a current instant in a status *Bad connection*. The conclusion about possible development of this attack is done if the connections quantity exceeds a preset threshold.

For construction of these detection rules an agent AD-P1 should use the reserved attribute N (number of the connections possessing defined parameters) of defined concept "current connection" (see the table A.1).

### Script A. The supervision of a quantity of the "bad" connections

A.1. The condition of the rule determines that the message about the connection transferred to a status *Bad connection* has come (see fig.A9, line 1 in the upper window).

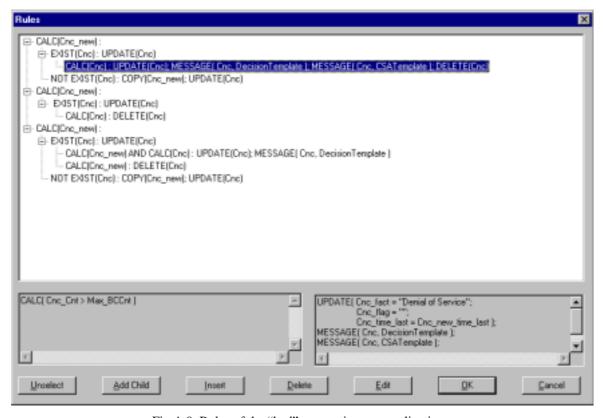


Fig.A.9. Rules of the "bad" connections generalization

- A.1.1. The rule determines that already there are connections with the same host in a status *Bad connection* and magnifies a value *N* on 1 in appropriate record of concept "current connection" (see Fig.A.9, line 2 in the upper window). For a search of the necessary record the following parameters are used: *IP address*, *Port destination* and a value *Bad connection* of the connection status in a field *Flag*.
  - A.1.1.1. If the value *N* has exceeded a preset maximum threshold, then the value *Denial of service* (*Syn Flood*) is assigned to a field *Fact*, and the instant, when the last connection has transferred to a status *Bad connection*, is recorded in a field *Time last tcp* (see Fig.A.9, line 3 in the upper window, and also contents of two windows below, this rule is marked). The message about the detected fact is transmitted to all agents *IDA1* of other hosts, and also to agent *IDA2* of the same defended host.
- A.1.2. The rule determines that for the present there are no connections with the host that send message in a status  $Bad\ connection$  (see Fig.A.9, line 4 in the upper window). In this case a new record with appropriate parameters from the obtained message is registered, and the value 1 is assigned to N.
- B.1. The rule determines that the message about the connection breaking in a status *Bad connection* on the life time has come (see Fig.A.9, line 5 in the upper window).
  - B.1.1. The rule discovers in the agent's database an appropriate record and reduces the value *N* on 1 in this record (see Fig.A.9, line 6 in the upper window).
    - B.1.1.1. If the value N became equal 0, then the record is deleted (see Fig.A.9, line 7 in the upper window).
- C.1. The rule determines that the message about the beginning of the new connection has come from some host (see Fig.A.9, line 8 in the upper window).
  - C.1.1. The rule discovers in the agent's database an appropriate record and magnifies the value *N* on 1 in this record (see Fig.A.9, line 9 in the upper window). For a search of the necessary record the following parameters are used: *IP address* and the value of the connection status *Half open connection*.
    - C.1.1.1. The rule checks that the open connections count time has not exceeded a preset time interval, and the open connections quantity has exceeded a preset threshold (see Fig.A.9, line 10 in the upper window). If this condition appears true, then a conclusion about port scanning is generated and the message about this fact is transferred to the agent *IDA2* of the same host.
    - C.1.1.2. The rule checks that the open connections count time has exceeded a preset time interval (see Fig.A.9, line 11 in the upper window). In this case the relevant record is deleted.
  - C.1.2. This rule works if the rule C.1.1 has not found a necessary record in the agent's database (see Fig.A.9, line 12 in the upper window). In this case a new record with parameters from the arrived message is registered, and the value 1 is assigned to N.

#### A1.2.3. Scripts of the agent *AD-P2* operation

The agent *AD-P2* is responsible for detection of attacks *Port scanning* (on an application layer), *Finger search* and *Buffer overflow* (Fig.A.10). Formal tag for these types attacks detection is a presence of a defined string in a data field of the input tcp-message addressed to the defined port. For example, under port scanning the tcp-packet field *Data* can include a substring "*expn cybercop*" at call to a port 25. The collection of examples of the similar sort possible combinations is established in the auxiliary table *PT*. One of three possible decisions is compared to each combination in this table:

• Port scanning,

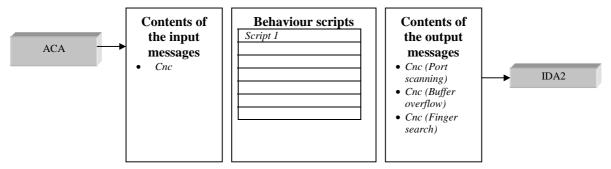


Fig.A.10. Common scheme of the agent *AD-P2* operation

- Finger search,
- Buffer overflow.

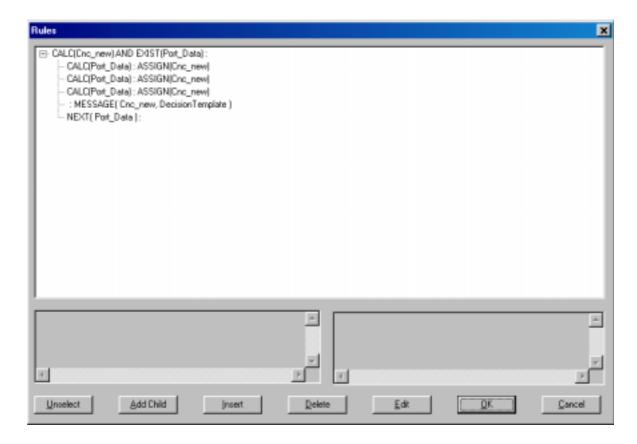


Fig.A.11. Rules of the agents AD-P2 behaviour

A.1. The rule uses data from a description of the current connection, to which the input tcp-message concerns (see Fig.A.11, line 1 in the upper window). At this *tcp*-message preprocessing stage the actual information (in particular a port number and string from a *tcp*-packet data field) was transferred to the connection description. The rule receives these data and checks them on a presence of an identical data in the auxiliary table *PT*.

A.1.1/2/3. If the rule A1 has detected one of three enumerated attack classes, then the next three subordinate rules register this result in a field Fact of the current connection description (see Fig.A.11, lines 2-4 in the upper window).

A.1.4. At detection of the attack fact this rule generates the appropriate message to the agent *IDA2* (see Fig. A.11, line 5 in the upper window).

# A1.2.4. Scripts of the agent AIA operation

The common scheme of the agent AIA operation is represented in Fig.A.12.

Contents of the input messages

• connect

Contents of the output messages

• connect

Contents of the output messages

• connection

Agent daemon

IDA2

Fig.A.12. Common scheme of the agent AIA operation

### Script A. Check of the user's name and password

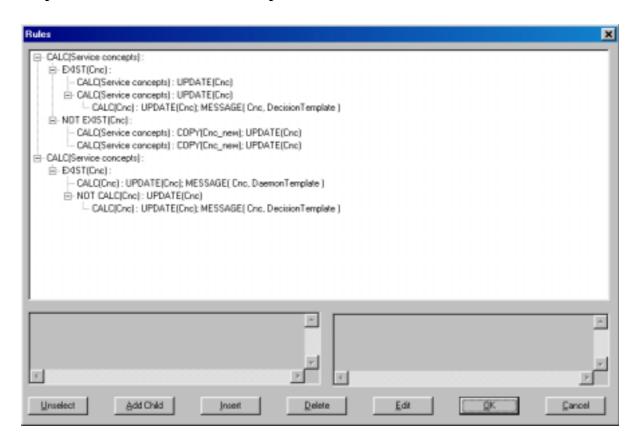


Fig.A.13. Rules of the agents AIA behaviour

The agents AIA behaviour is determined by two generalized rules A.1 and B.1. The first rule makes an identification of a login name, and the second rule checks a password.

A.1. The rule determines that an input message contains a login name (see a Fig.A.13, line 1).

- A.1.1. The rule determines that an attempt of a name input for the connection establishment is not the first within the current connection (see Fig.A.13, line 2).
  - A.1.1.1. The rule determines that the input name is contained in the table of names of the users having the right on operation with a host (see a Fig.A.13, line 3). In this case the value *user\_ok* is recorded in a field *Fact* of the current connection description.
  - A.1.1.2. The rule determines that the table of names of the users does not contain the input name (see Fig.A.13, line 4). In this case a value of the access attempts counter is magnified on 1.
    - A.1.1.2.1. If the access attempts quantity is more than a preset threshold, then a login name guessing conclusion is adduced, and a message about this fact is sent to the agent *IDA2* (see Fig.A.13, line 5).
- A.1.2. The rule determines that the login name input attempt for a connection establishment is the first within the current connection (see Fig.A.13, line 6).
  - A.1.2.1. The rule determines that the input name is contained in the table of names of the users having the right on operation with a host (see a Fig.A.13, line 7). In this case the value *user\_ok* is recorded in a field *Fact* of the current connection description. The value of the host access attempts counter remains equal 0.
  - A.1.2.2. The rule determines that the table of names of the users does not contain the input name (see Fig.A.13, line 8). In this case the value 1 is assigned to the access attempts counter.
- B.1. The rule determines that a user's password is in the input message data field (see a Fig.A.13, line 9).
  - B.1.1. The rule checks that a login name existing in the users' names and passwords table was entered (see a Fig.A.13, line 10). This table contains the names and passwords of the users having the right on operation with a host.
    - B.1.1.1. The rule determines that the input password is contained in the users' names and passwords table and corresponds to the login name input earlier (see Fig.A.13, line 11). The user's access rights are assigned on the basis of the same table. The message about a successful access with the indication of access rights in the current connection description field *Access* is transmitted to the agent *AD-E*.
    - B.1.1.2. The rule determines that the input password is incorrect and magnifies the access attempts counter on 1 (see Fig.A.13, line 12).
      - B.1.1.2.1. If the host access attempts quantity is more than a preset threshold, then a login name and password guessing conclusion is adduced, and a message about this fact is sent to the agent *IDA2* (see Fig.A.13, line 13).

## A1.2.5. Scripts of the agent ACA operation

The common scheme of the agent ACA operation is represented in Fig.A.14.

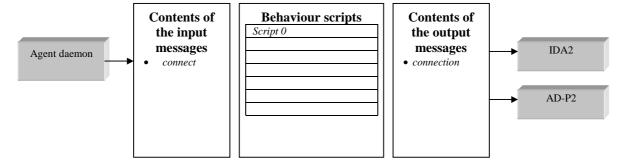


Fig.A.14 Common scheme of the agent ACA operation

### Script A. Check of the access rights

The message including the user's access rights and the input *tcp*-message data field contents arrives to the agent *ACA* (see Fig.A.15, line 1). If the actions defined in the data field correspond to the available rights, then the first rule transmits the obtained message to the agent *AD-P2* (see Fig.A.15, line 2). If the access violation attempt is detected, then the second rule transmits this fact to the agent *IDA2* (see Fig.A.15, line 3, and also contents of two windows below, this rule is marked).

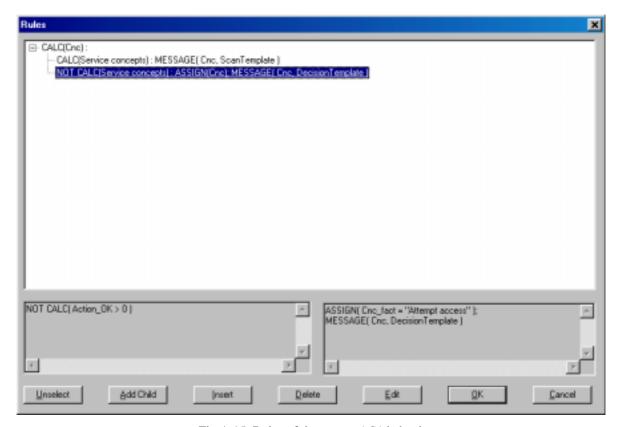


Fig.A.15. Rules of the agents ACA behavior

### A1.2.6. Scripts of the agent *IDA1* operation

The common scheme of the agent *IDA1* operation is represented in Fig.A.16.

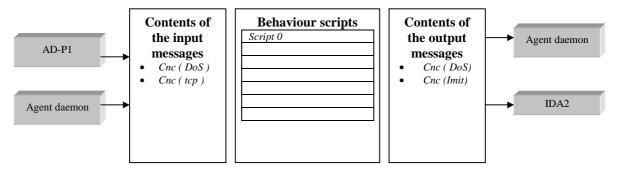
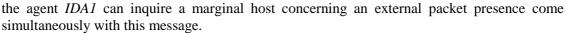


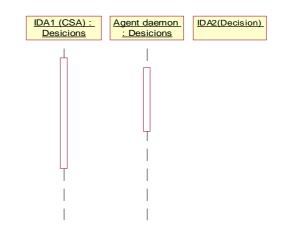
Fig.A.16. Common scheme of the agent *IDA1* operation

The agents *IDA1* fulfill the task of the *Combined spoofing attack* detection. In Case study the base mechanisms of this task decision are realized. The sequence diagram (Fig.A.17) explains an

order of the agents' interaction. For implementation of a more complete task solution the following agents' dialogues is necessary to add in the scheme of the agents' interaction.

- At obtaining a message about the fact of the presumable attack *Denial of service* from a host H, the agent *IDA1* can inquire this host for a supposition confirmation.
- At obtaining the tcp-message with a flag SYN from the host undergone a *Denial of service*,





## Script A. Combined spoofing attack detection

The rules of the agent *IDA1* behaviour for the Combined spoofing attack detection are represented in Fig.A18.

A1. The rule determines that the message on the possible attack *Denial of service* is obtained (see Fig.A18, line 1). On the basis of the obtained data the rule transfers to the agent *AD-E* a message with the contents represented in the left part of the table 2 (from the description of the agents *AD-E* behaviour).

A2. This rule fixes the *tcp*-message with a flag SYN obtained from the attacked host (see Fig.A18, line 2).

A3. This rule fixes a termination of the "hand shake" on behalf of the host undergone *attack Denial* of service, and concludes a fact about an imitation of the connection on behalf of this host (see Fig.A18, line 3, and also contents of two windows below, this rule is marked). The message on the detected fact is transmitted to the agent *IDA2*.

### A1.2.7. Scripts of the agent IDA2 operation

The common scheme of the agent *IDA1* operation is represented in Fig.A.19.

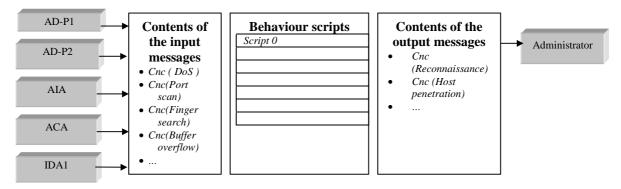


Fig.A.19. Common scheme of the agent IDA2 operation

A.1. The rule is fulfilled, if the message about a host access attempt by a login name and password guessing or a *Combined spoofing attack* detection is arrived (see Fig.A.20, lines 1 and 2). In this case the rule parses a quantity of the similar sort facts for a preset time period. If this quantity exceeds a preset threshold, then a conclusion about a *Host penetration* is generated and the warning is dispatched to the administrator.

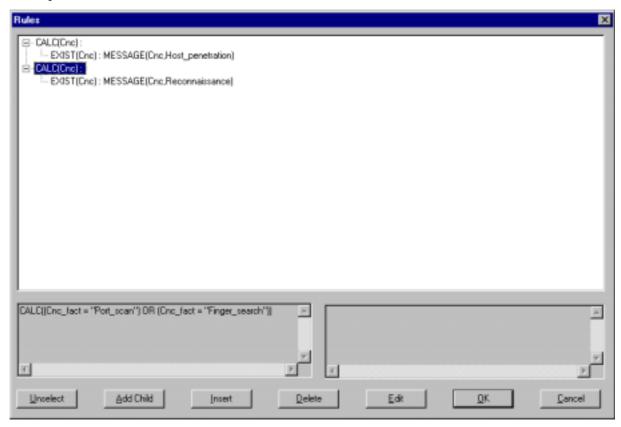


Fig.A20. Rules of the agents IDA2 behaviour

AA.2. The rule is fulfilled, if the message about the *Port scanning* or *Finger search* facts has arrived (see Fig.A.20, lines 3 and 4, and also contents of two windows below, this rule is marked). In this case the rule parses a quantity of the similar sort facts for a preset time period. If this quantity exceeds a preset threshold, then a conclusion about a *Reconnaissance* is generated and the warning is dispatched to the administrator.